

# Ptolemy

Balances expressiveness and modular reasoning  
for aspect-oriented software development.

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# Tutorial Outline

## ❖ Why Ptolemy? What problems does it solve?

### ❖ Two precursors

➤ Implicit Invocation and Aspect-orientation

## ❖ Ptolemy and how it solves these problems.

## ❖ Main Language Features

❖ Declarative, typed events (join points in AO terms)

❖ Declarative, typed event announcement (no AO term)

❖ Declarative, typed event registration (advising in AO terms)

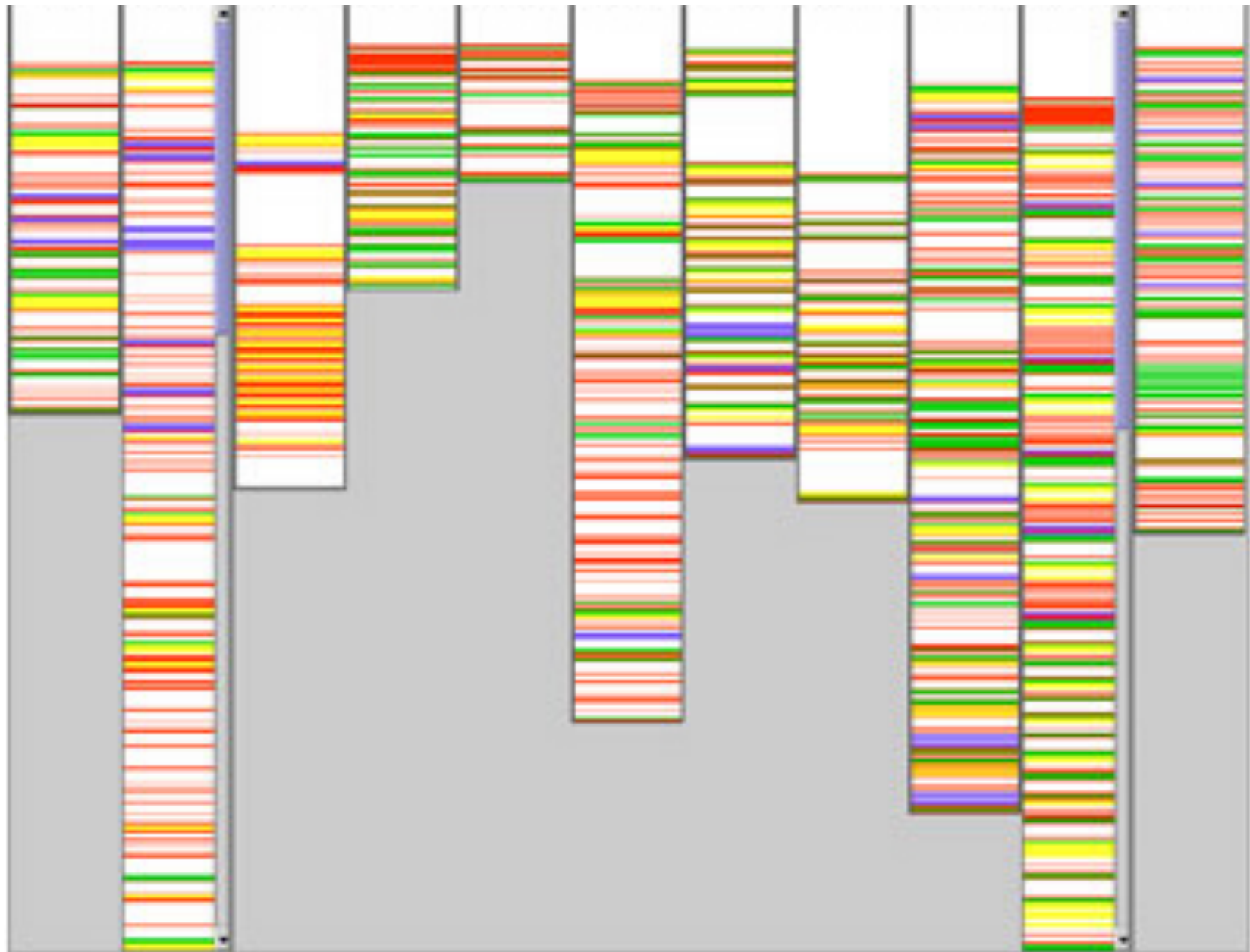
❖ Quantification based on event types (same as the AO term)

# Tutorial Outline

- ❖ Modular Verification Features
  - ❖ Translucid Contracts (no AO term)
- ❖ Where to use Ptolemy Features?
  - ❖ vs. Aspect-orientation,
  - ❖ vs. Implicit Invocation
- ❖ State of Tools
- ❖ Opportunities to Contribute
- ❖ Conclusion

One shall not have to choose between reasoning and separation.

**WHY PTOLEMY?**



Color-coded representation of about 19K LOC at ASML: different colors represent different concerns in the system.  
Courtesy: Bruntink, Deursen and Tourwé

# Need for Improved Separation

- ❖ Some concerns hard to modularize
- ❖ Number of proposals: Units [Flatt and Felleisen], Mixin [Bracha and Cook], Open Classes [Clifton et al.], Roles [Kristensen and Osterbye], Traits [Scharli et al.], Implicit Invocation [Garlan, Notkin, Sullivan et al.], Hyperslices [Ossher and Tarr], Aspects [Kiczales et al.], etc
- ❖ Shows that there is a real need

# Two similar ideas

- ❖ Implicit invocation (II) vs. Aspect-orientation (AO)
- ❖ ... both effective for separation of concerns
- ❖ ... both criticized for making reasoning hard
  - ❖ II criticized in early/late 90's
  - ❖ AO criticized in early 2000's
- ❖ Ptolemy is designed to
  - ❖ combine best ideas from II and AO
  - ❖ ... and to make reasoning easier

[JHotDraw – Gamma et al.]

# **RUNNING EXAMPLE**



# Elements of a Drawing Editor

- ❖ Elements of drawing

  - ❖ Points, Lines, etc

  - ❖ All such elements are of type Fig

- ❖ Challenge I: Modularize display update policy

  - ❖ Whenever an element of drawing changes — Update the display

- ❖ Challenge II: Impose application-wide restriction

  - ❖ No element may move up by more than 100

# Figure Elements

```
1 abstract class Fig {  
2 }
```

- ❖ Fig – super type for all figure elements
  - ❖ e.g. points, lines, squares, triangles, circles, etc.

# Point and its Two Events

```
1. class Point extends Fig {
2     int x;
3     int y;
4     void setX(int x) {
5         this.x = x;
6     }
7     ..
8     void makeEqual(Point other) {
9         if(!other.equals(this)) {
10             other.x = this.x;
11             other.y = this.y;
12         }}}
```

- ❖ Changing Fig is different for two cases.
- ❖ Actual abstract event inside makeEqual is the true branch.

Reiss'92, Garlan and Notkin'92

# IMPLICIT INVOCATION

# Key Ideas in II

- ❖ Allow management of name dependence
  - ❖ when “Point’s coordinates changes” update Display
  - ❖ ... but Point shouldn’t depend on Display
  - ❖ ... complicates compilation, test, use, etc
- ❖ Components (subjects) declare events
  - ❖ e.g. when “Point’s coordinates changes”
  - ❖ provide mechanisms for registration
  - ❖ ... and for announcement
- ❖ Components (observers) register with events
  - ❖ e.g. invoke me when “Point’s coordinates changes”
- ❖ Subjects announce events
  - ❖ e.g. when “Point’s coordinates changes”
  - ❖ “change in coordinates” event announced

## II: Components Declare Events

```
1 abstract class Fig {
2     List changeObservers;
3     void announceChangeEvent(Fig changedFE) {
4         for(ChangeObserver o : changeObservers) {
5             o.notify(changedFE);
6         }
7     }
8     void registerWithChangeEvent(ChangeObserver o) {
9         changeObservers.add(o);
10    }
11 }
12 abstract class ChangeObserver {
13     void notify(Fig changedFE);
14 }
```

## II: Components Announce Events

```
1 class Point extends Fig {
2   int x; int y;
3   void setX(int x) {
4     this.x = x;
5     announceChangeEvent(this);
6   }
7   void makeEqual(Point other) {
8     other.x = this.x; other.y = this.y;
9     announceChangeEvent(other);
10  }
11 }
```

- ❖ Event announcement explicit, helps in understanding
- ❖ Event announcement flexible, can expose arbitrary points

## II: Component Register With Events

```
1 class Update extends ChangeObserver {
2     Fig last;
3     void registerWith(Fig fe) {
4         fe.registerWithChangeEvent(this);
5     }
6     void notify(Fig changedFE){
7         this.last = changedFE;
8         Display.update();
9     }
10 }
```

- ❖ Registration explicit and dynamic, gives flexibility
- ❖ Generally deregistration is also available



## II: Disadvantages

- ❖ Coupling of observers to subjects

```
void registerWith(Fig fe) {  
    fe.registerWithChangeEvent(this); ...  
}
```

- ❖ Lack of quantification

```
void registerWith(Point p) {  
    p.registerWithChangeEvent(this);  
}  
  
void registerWith(Line l) {  
    l.registerWithChangeEvent(this);  
}
```

## II: Disadvantages

❖ No ability to replace event code

```
class MoveUpCheck extends ... {  
    void notify(Fig targetFE, int y, int delta) {  
        if (delta < 100) { return targetFE }  
        else{throw new IllegalArgumentException()}  
    }  
}
```

Kiczales et al. 97, Kiczales et al. 2001

# ASPECT-BASED SOLUTIONS

# Key Similarities/Differences with II

- ❖ Events  $\equiv$  “join points”
  - ❖ AO: pre-defined by the language/ II: programmer
  - ❖ AO: Implicit announcement/ II: explicit
- ❖ Registration  $\equiv$  Pointcut descriptions (PCDs)
  - ❖ AO: declarative
- ❖ Handlers  $\equiv$  “advice” register with sets of events
- ❖ Quantification: using PCDs to register a handler with an entire set of events

# Aspect-based Solution

```
1  aspect Update {
2  Fig around(Fig fe) :
3      call(Fig+.set*(..)) && target(fe)
4      || call(Fig+.makeEq*(..)) && args(fe){
5      Fig res = proceed(fe);
6      Display.update();
7      return res;
8}
```

# Advantages over II

- ❖ Ease of use due to quantification
- ❖ By not referring to the names, handler code remains syntactically independent

# Limitations: Fragility & Quantification

- ❖ Fragile Pointcuts: consider method “settled”

```
1 Fig around(Fig fe) :  
2 call(Fig+.set*(..)) && target(fe)  
3 || call(Fig+.makeEq*(..)) && args(fe){  
4 ...
```

- ❖ Quantification Failure: Arbitrary events not available

```
1 Fig setX(int x){  
2   if (x.eq(this.x)) { return this; }  
3   /* abstract event change */  
4   else { this.x = x; return this; }  
5 }
```

# Limitations: Context access

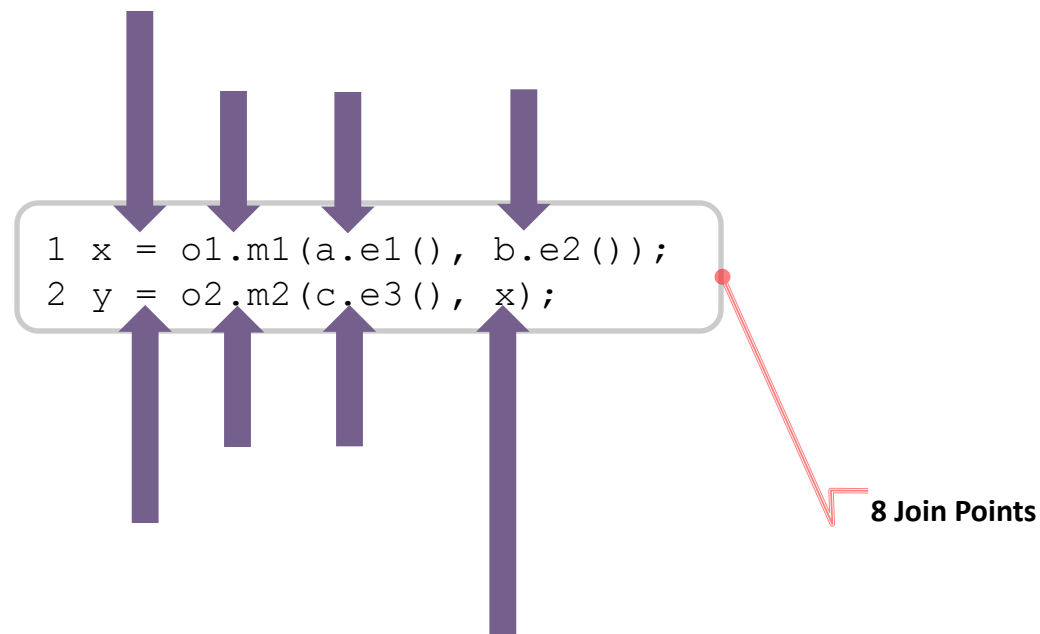
## ❖ Limited Access to Context Information

- ❖ Limited reflective interface (e.g. “thisJoinPoint” in AJ)
- ❖ Limited Access to Non-uniform Context Information

```
1 Fig around(Fig fe) :  
2 call(Fig+.set*(..)) && target(fe)  
3 || call(Fig+.makeEq*(..)) && args(fe){  
4 ...
```



# Limitations: Pervasive Join Point Shadows



- ❖ For each join point shadow, all applicable aspect should be considered (whole-program analysis)

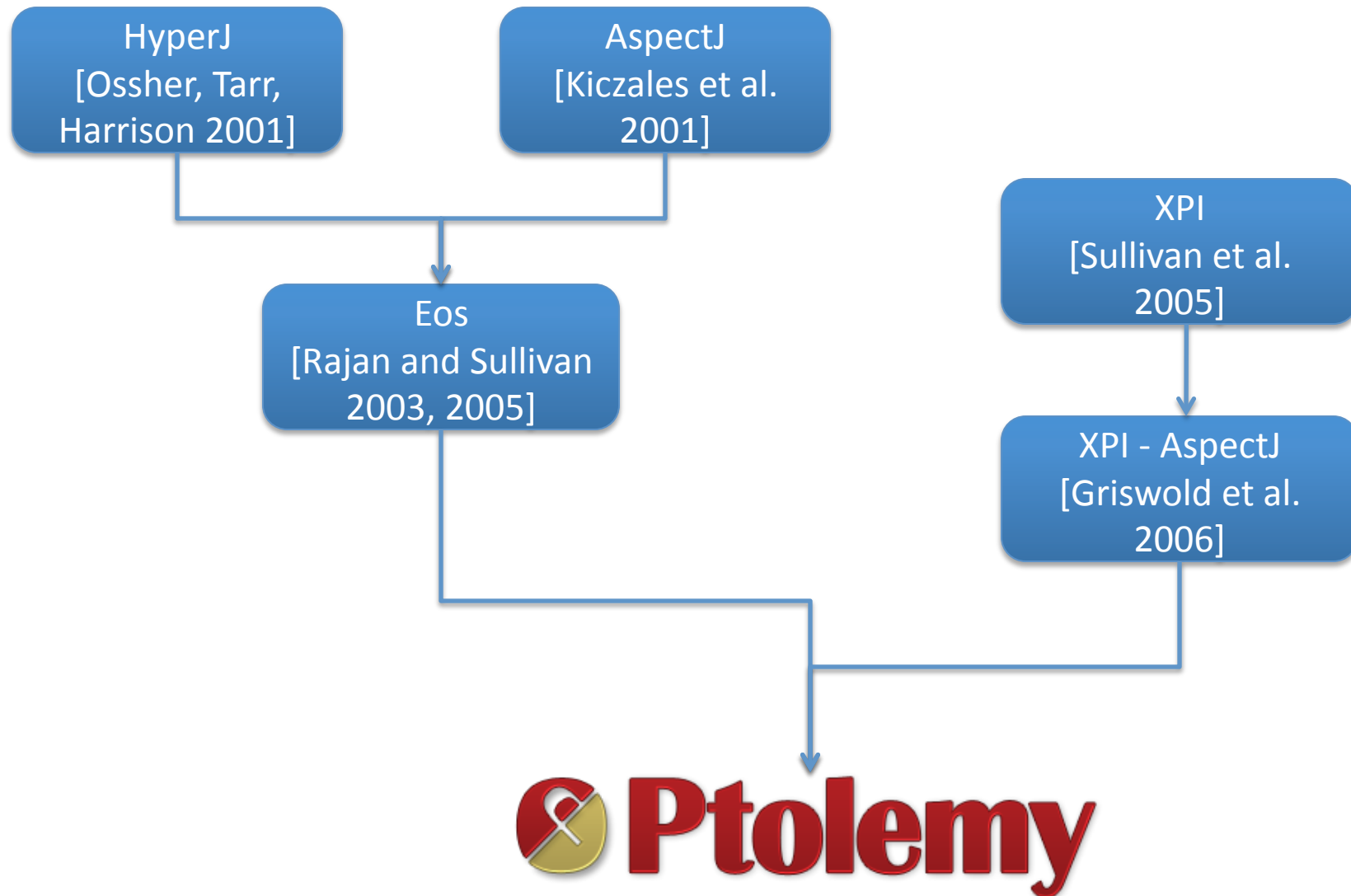


Ptolemy (Claudius Ptolemaeus), fl. 2d cent. A.D., celebrated Greco-Egyptian mathematician, astronomer, and geographer.

 **Ptolemy**

[ptolemy.cs.iastate.edu](http://ptolemy.cs.iastate.edu)

# Evolution of the Ptolemy Language



# Design Goals of Ptolemy

- ❖ Enable modularization of crosscutting concerns, while preserving encapsulation of object-oriented code,
- ❖ enable well-defined interfaces between object-oriented code and crosscutting code, and
- ❖ enable separate type-checking, separate compilation, and modular reasoning of both OO and crosscutting code.

# First and foremost

- ❖ Main feature is event type declaration.
- ❖ Event type declaration design similar to API design.
  - ❖ What are the important abstract events in my application?
  - ❖ When should such events occur?
  - ❖ What info. must be available when such events occur?
- ❖ Once you have done it, write an event type declaration.

# Declaring an Event Type

```
Fig event Changed {  
    Fig fe;  
}
```



**Event Type  
Declaration**

# Declaring an Event Type

```
Fig event Changed {  
    Fig fe;  
}
```



**Event Type  
Declaration**

- ❖ Event type is an abstraction.
- ❖ Declares context available at the concrete events.
- ❖ Interface, so allows design by contract (DBC) methodology.

# Announcing Events in Ptolemy

## Subject

```
1 class Fig {bool isFixed;}
2 class Point extends Fig{
3   int x, y;
4   Fig setX(int x){
5     announce Changed(this) {
6       this.x = x; return this;
7     }
8   }
9 }
```

Event  
Announcement

❖ Explicit, more declarative, typed event announcement.



# More Event Announcements

## Subject

```
class Point extends Fig{  
  ..  
  Fig moveUp(int delta){  
    announce MoveUpEvent(this){  
      this.y += delta; return this;  
    }  
  }  
}
```

Event  
Announcement

❖ Explicit, more declarative, typed event announcement.



# Quantification Using Binding Decls.

## ❖ Binding declarations

❖ Separate “what” from “when” [Eos 2003]

### Observer (Handler)

```
class DisplayUpdate {
```

```
  when Changed do update;
```

```
}
```

Quantification

# Dynamic Registration

- ❖ Allow dynamic registration
  - ❖ Other models can be programmed

## Observer (Handler)

```
class DisplayUpdate {  
  
    void DisplayUpdate() { register(this) }  
  
    Fig update(Changed next) {  
  
    }  
  
    when Changed do update;  
}
```

Registration

Quantification

# Controlling Overriding

- ❖ Use `invoke` to run the continuation of event
- ❖ Allows overriding similar to AspectJ

## Observer (Handler)

```
class DisplayUpdate {
```

```
void DisplayUpdate() { register(this);
```

Registration

```
Fig update(Changed next) {
```

```
System.out.println("Before Invoke");
```

```
next.invoke();
```

```
Display.update();
```

```
System.out.println("After Invoke");
```

```
}
```

Running  
continuation of  
the event

```
when Changed do update;
```

```
}
```

Quantification

# Exercise 0: Get the distribution

❖ Go to the URL to download Ptolemy1.2 Beta5  
<http://ptolemy.cs.iastate.edu/ptolemy-ase11.zip>

to download the zip file ***ptolemy-ase-11.zip***

❖ Unzip the contents at a convenient location, while preserving its directory structure

❖ Start Eclipse and change your workspace to

`$Download_Location$/pyc/workspace`

# Exercise 1: Figure Editor Example

- ❖ Browse code in 00-FigureExample Project
- ❖ [a] Open file FEChanged.java
  - ❖ Note return type and context variables of event declaration FEChanged
- ❖ [b] Open file Point.java
  - ❖ Note event announcements in `setX`, `setY`, `moveBy`
  - ❖ Note different context in method `makeEqual`.
    - ❖ Everywhere else context variable `changedFE` is bound to **this**, but in this method it is bound to `other`.

# Exercise 1: Figure Editor Example

- ❖ [c] Open file `DisplayUpdate.java`

- ❖ Note the binding declaration

- `when FEChanged do update`

- ❖ Note the **register** statements

- `register (..)`

- It registers the receiver object to listen to events mentioned in the binding declarations



# EXTENDED EXAMPLE

# Expressions and Operations

ASTNodes

$e ::= v$

| (lambda (v) . e )

| (e e)

$e ::= \dots$

| true | false | Num

| e == e | e <= e

| e && e | e `||` e

| e + e | e \* e | e - e

Eval

$E: e \implies e'$

Checker

$T \vdash e : t$

Printer

# Goal: Separation of Concerns

## ASTNodes

```
e ::= v
    | (lambda (v) . e)
    | (e e)

e ::= ...
    | true | false | Num
    | e == e | e <= e
    | e && e | e `||` e
    | e + e | e * e | e - e
```

## Eval

$E: e \Rightarrow e'$

## Checker

$T \vdash e : t$

## Printer

# Goal: Separation of Operations

## AST Nodes

```
e ::= v
    | (lambda (v) . e)
    | (e e)

e ::= ...
    | true | false | Num
    | e == e | e <= e
    | e && e | e `||` e
    | e + e | e * e | e - e
```

## AST Events

### Eval

$E: e \implies e'$

### Checker

$T \dashv\vdash e : t$

### Printer

Enabling modular verification

# CONTRACTS IN PTOLEMY

**DEMO**

# Conclusion

- ❖ Motivation: intellectual control on complexity essential
  - ❖ Implicit invocation (II) and aspect-orientation (AO) help
  - ❖ ... but have limitations
- ❖ Ptolemy: combine best ideas of II and AO
  - ❖ Quantified, typed events + arbitrary expressions as explicit events
  - ❖ Translucid contracts
- ❖ Benefits over implicit invocation
  - ❖ decouples observers from subjects
  - ❖ ability to replace events powerful
- ❖ Benefits over aspect-based models
  - ❖ preserves encapsulation of code that signals events
  - ❖ uniform and regular access to event context
  - ❖ robust quantification
- ❖ Last but not least, more modular reasoning

# Opportunities to Contribute

- ❖ Language design efforts
  - ❖ **Ptolemy# to come out in June, testing underway (Extension of C#)**
  - ❖ Transition to less front-end changes (for PtolemyJ)
- ❖ Verification efforts
  - ❖ More expressive support for embedded contracts
  - ❖ Practical reasoning approaches for heap effects
  - ❖ Better verification error reporting

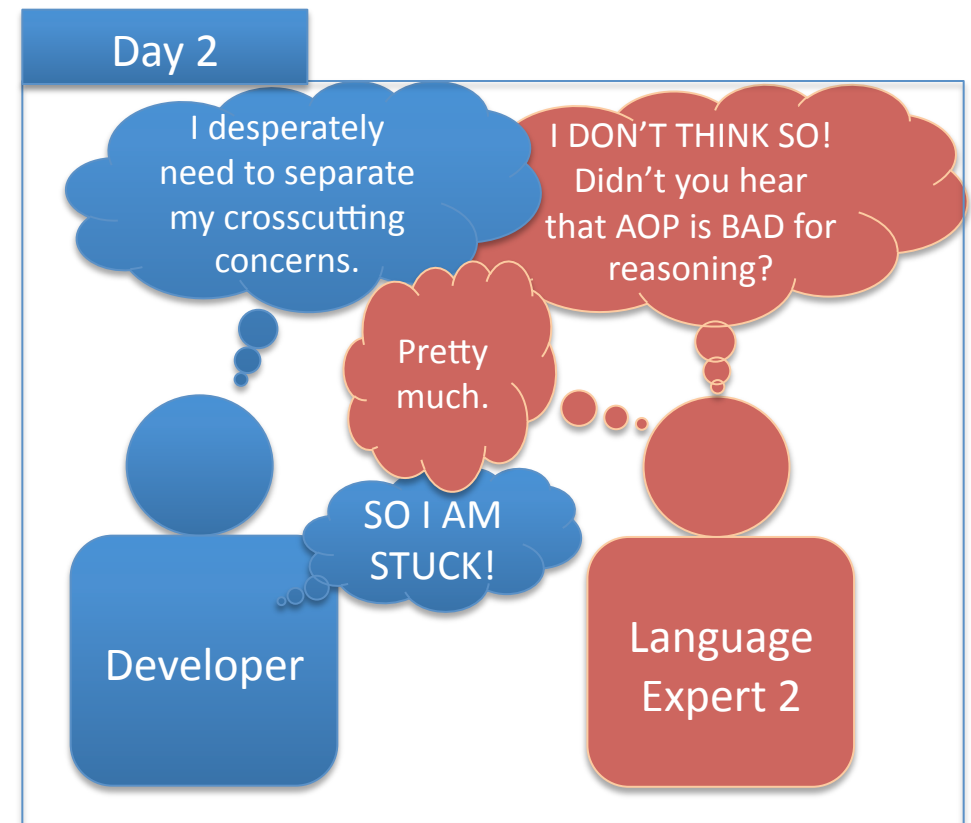


# Opportunities to Contribute

- ❖ Case study efforts – compiler supports metrics
  - ❖ Showcase applications, examples for Ptolemy
  - ❖ Comparison with other languages/approaches
- ❖ Infrastructure efforts
  - ❖ Support in Eclipse, other IDEs
  - ❖ Better error reporting, recovery
- ❖ Language manuals, descriptions,...

**All are welcome!!!**

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