#### Software Development for Embedded Systems

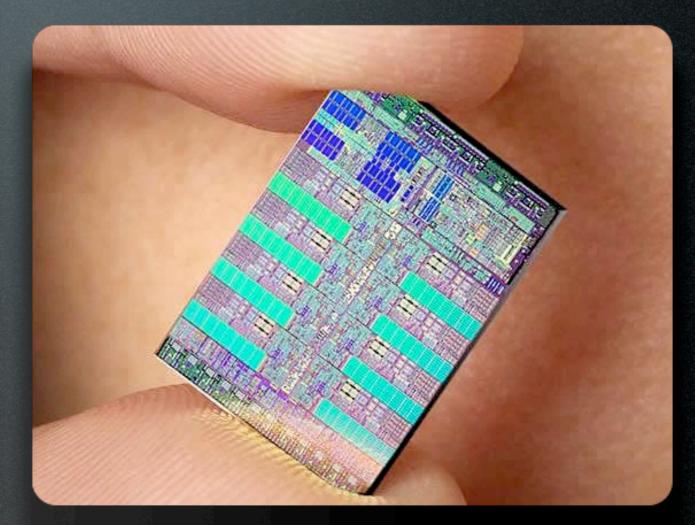
Witawas Srisa-an CSCE 496: Embedded Systems Design and Implementation

#### Reminder

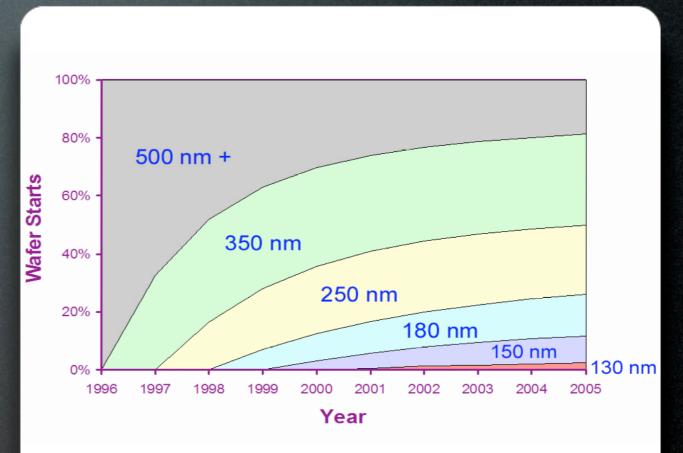
- Homework on Gaussian Filter is due on Wednesday before 11:59 pm
  - if you are having difficulties reading the provided image file, be sure to read the help page on fopen in the Stretch IDE
- Any questions about malloc?

- Proliferation of complex embedded systems
  - powerful processors
    - feature-rich (e.g. Cell, Stretch, ARM)
    - advanced runtime support
      - similar features

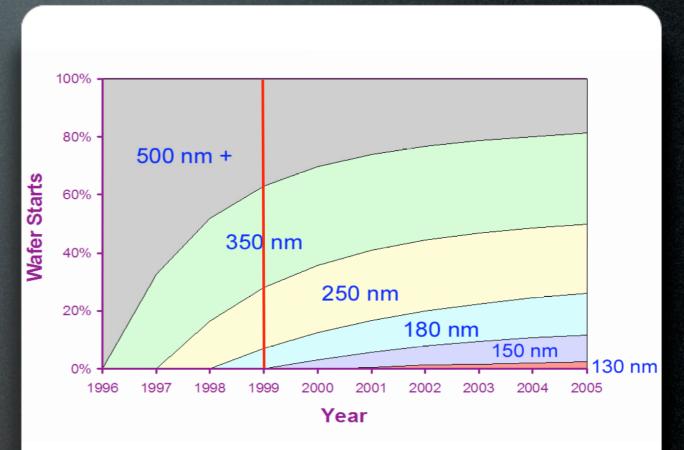
         found in processors
         for desktop and
         server systems (e.g.
         MMU, multi-core,
         fast bus, etc.)



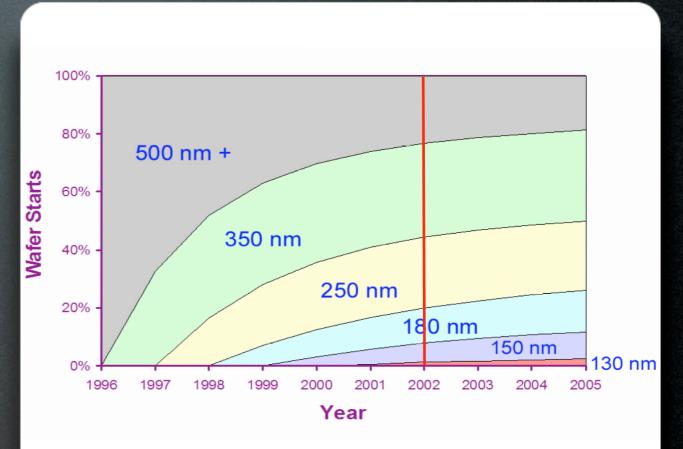
- We are reaping the major benefit of Moore's law
  - old processors don't go away
    - powerful enough for day-to-day applications
    - very low cost



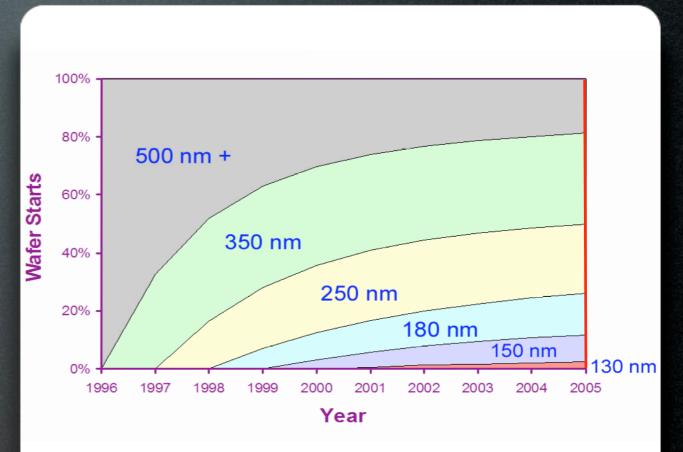
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- Complex software systems
  - why not? The hardware can handle it!
  - Avionic Software for Boeing ScanEagle UAV > 300,000 line of code



- Feature-rich runtime support
  - full-fledged desktop/ server operating systems in embedded devices
  - capability to run more complex software systems on these machines



- Leverage mobile/ embedded devices to provide services
  - <u>Telesensing</u> from Lucent
    - prevent Sudden Infant Death Syndrome, detect sleep apnea, etc.
  - Full-fledged web browsers, calendars, file servers, etc.



- Assume heterogeneous platforms
  - the billionth handset shipped in 2006
    - build for portability
    - build for generic input/output devices
  - about 20+ operating systems for these devices
    - build to interface with underlying runtime support features

#### • Summary

- large software systems now and larger in the future
  - more software reuse?
- assume heterogeneous platforms
  - must be portable
- providing similar runtime features to much more complex systems

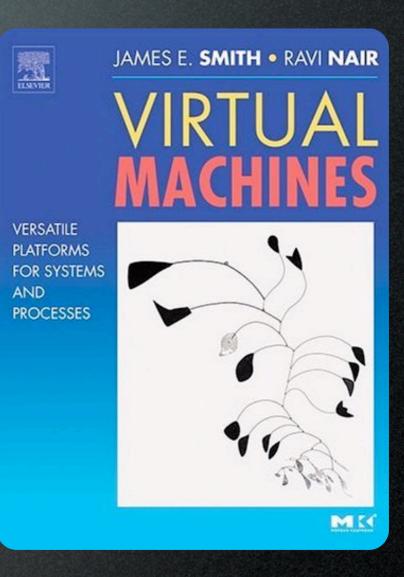
# Enabling Technologies

- Modern programming languages (e.g. Java, C#, VB.NET)
  - Object-oriented paradigm
    - promote code reuse
  - Virtual Machine (VM) based systems
    - achieve portability but require complex runtime support
      - now available in many embedded devices

# Virtual Machines

"A **virtual machine** is <u>software</u> that creates a <u>virtualized environment</u> between the <u>computer</u> <u>platform</u> and its operating system, so that the <u>end</u> <u>user</u> can operate software on an abstract machine."

Wikipedia

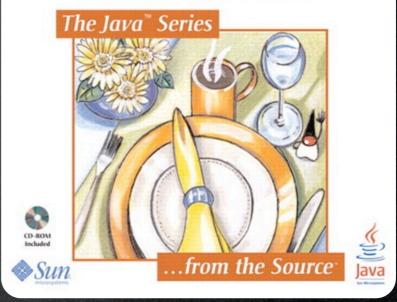


"A Java Virtual Machine (JVM) is virtual machine that <u>interprets</u> and <u>executes</u> Java bytecode. This code is most often generated by <u>Java</u> <u>language compilers...</u>"

Wikipedia

James Gosling • Bill Joy • Guy Steele • Gilad Bracha 🔸

#### The Java<sup>®</sup> Language Specification, Third Edition



- We'll look at the ones developed mainly by Sun Microsystems
  - HotSpot
  - CLDC HotSpot
  - KVM (Kilo Virtual Machine)
    - second most used VM in the world

developed by Sun?

• We'll look at the ones developed mainly by Sun Microsystems Trivia: What is the most used VM

CLDC HotSpot
KVM (Kilo Virtual Machine)
second most used VM in the world

#### • KVM vs. HotSpot

- KVM is interpretation based
  - maximum portability
- HotSpot combines interpretation and dynamic compilation
  - platform dependent

# Sidebar: Dynamic Compilation

- Interpreter is a big while loop with many case statements
  - each bytecode is translated to a predefine C/C+ function (e.g. new operator)
- Dynamic compiler takes each method and generates native code
  - can be optimized or non-optimized

# Sidebar: Dynamic Compilation

#### • Compilation strategies

- always compile (e.g. .NET Compact Framework, Jikes RVM)
- only compile frequently used methods
- Code size
  - a compiled method can be 6 to 8 times larger than its bytecode representation

# Sidebar: Dynamic Compilation

#### • Storage

- these compiled methods are stored in a dynamic memory region
  - separate code-cache or intermingled with objects in the heap
- Management strategies
  - flush when full, GC, etc.

- KVM vs. HotSpot
  - KVM uses simple mark-sweep-compact garbage collection
    - simple but long execution pauses
  - HotSpot uses generational garbage collection
    - more complex with higher runtime overhead, but shorter pauses

**"Garbage Collection (GC)** is a form of automatic <u>memory management</u>. The garbage collector attempts to reclaim <u>garbage</u>, or memory used by <u>objects</u> that will never again be accessed or mutated by the <u>application</u>."



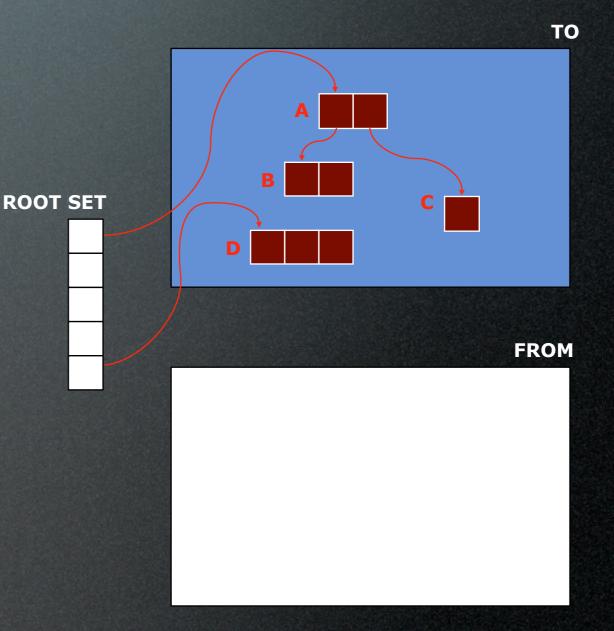
Algorithms for Automatic Dynamic Memory Management

Richard Jones Rafael Lins

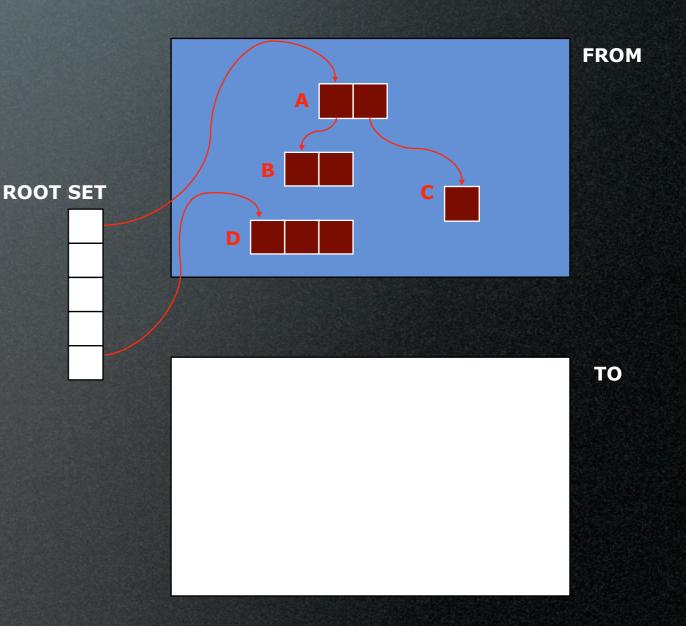


#### Wikipedia

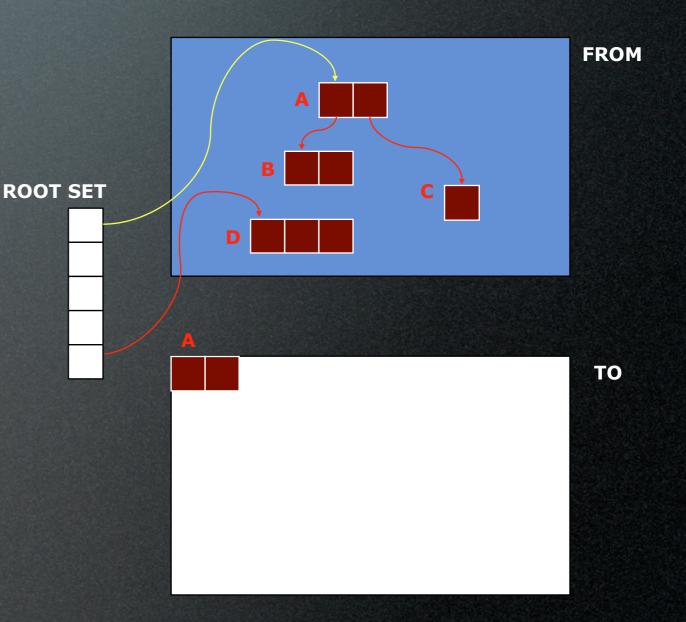
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- when the half is full, migrate surviving objects to the other half then allocate new objects from there



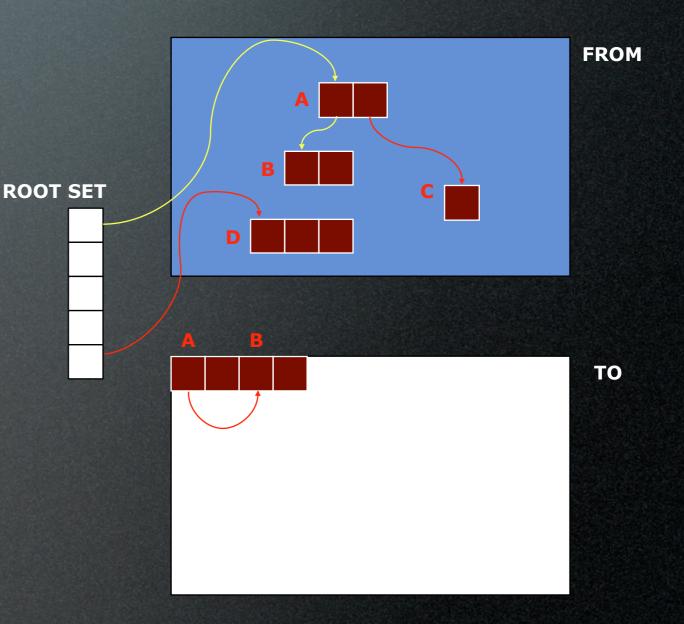
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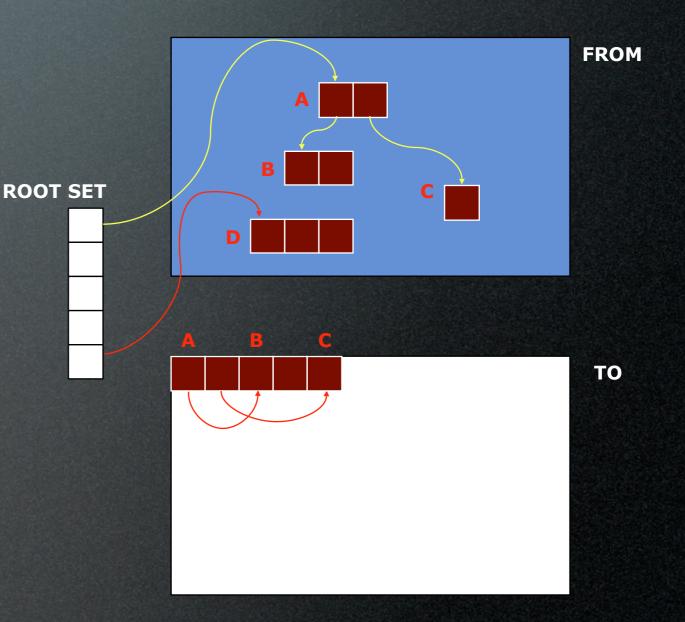
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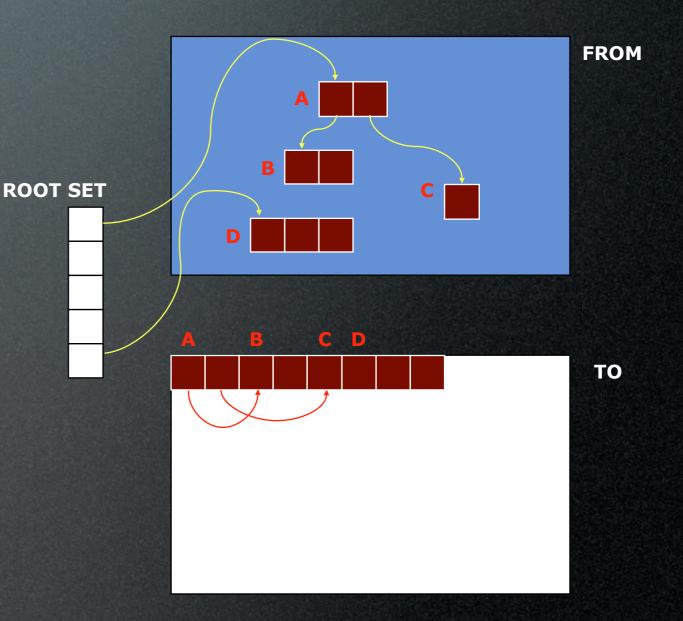
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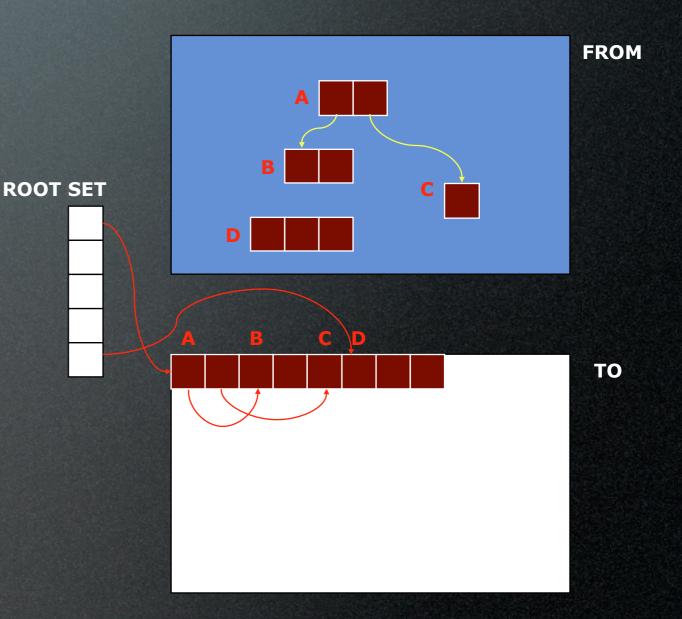
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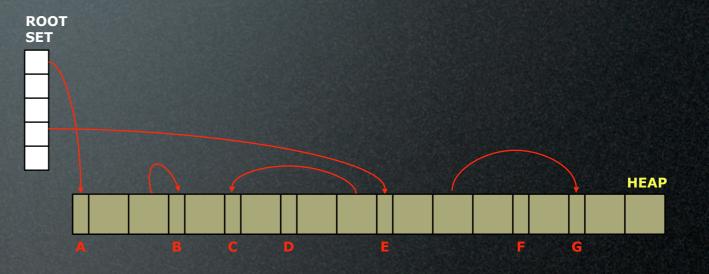
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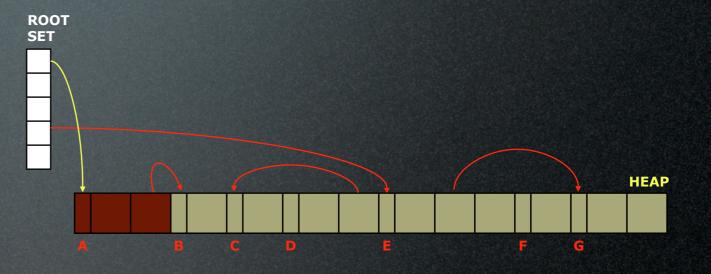
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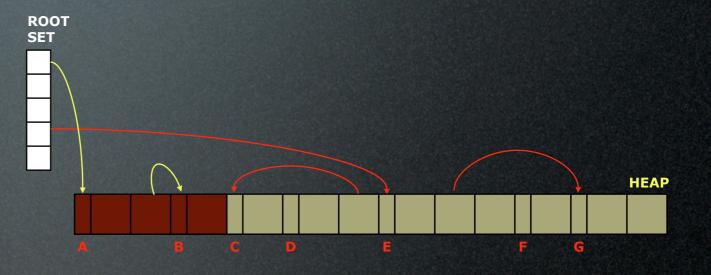
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  - collect the entire heap each time
  - when the heap is full, identify live objects (marking), then free dead objects (sweeping)



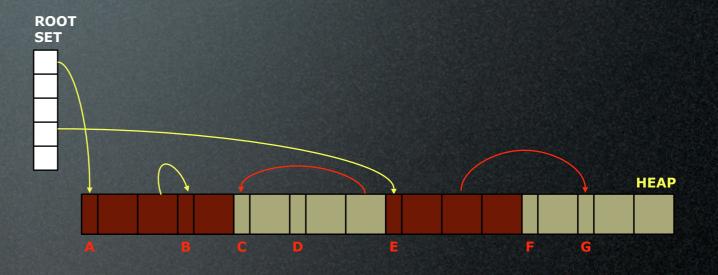
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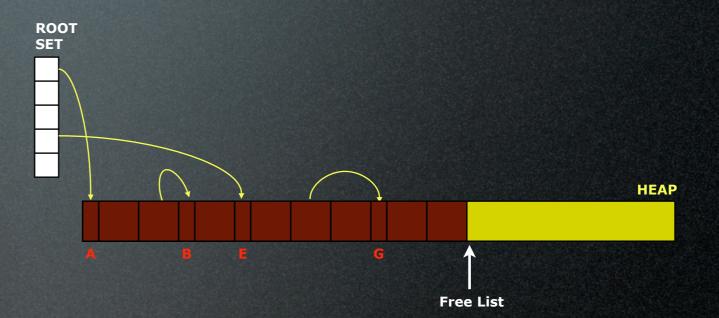
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- Mark-sweep-compact collector
  - collect the entire heap each time
  - similar to marksweep collector except that heap is compacted after sweeping



- But typical JVMs are not ready for deployment in realtime embedded systems
  - lacking real-time support
    - unpredictable execution of operations
    - no support for real-time threads
    - no priority inversion avoidance
    - unbounded garbage collection

• But typical JVMs are not ready for deployment in realtime embedded systems

# These topics will be the focus of the next few lectures

unpredictable execution of operations

• no support for real-time threads

no priority inversion avoidance

unbounded garbage collection