

Game-Theoretic View to Multiagent Systems

(Based on Shoham and Leyton-Brown (2008). *Multiagent Systems: Algorithmic, Game-Theoretic, and Logical Foundations*, Cambridge.)

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Introduction

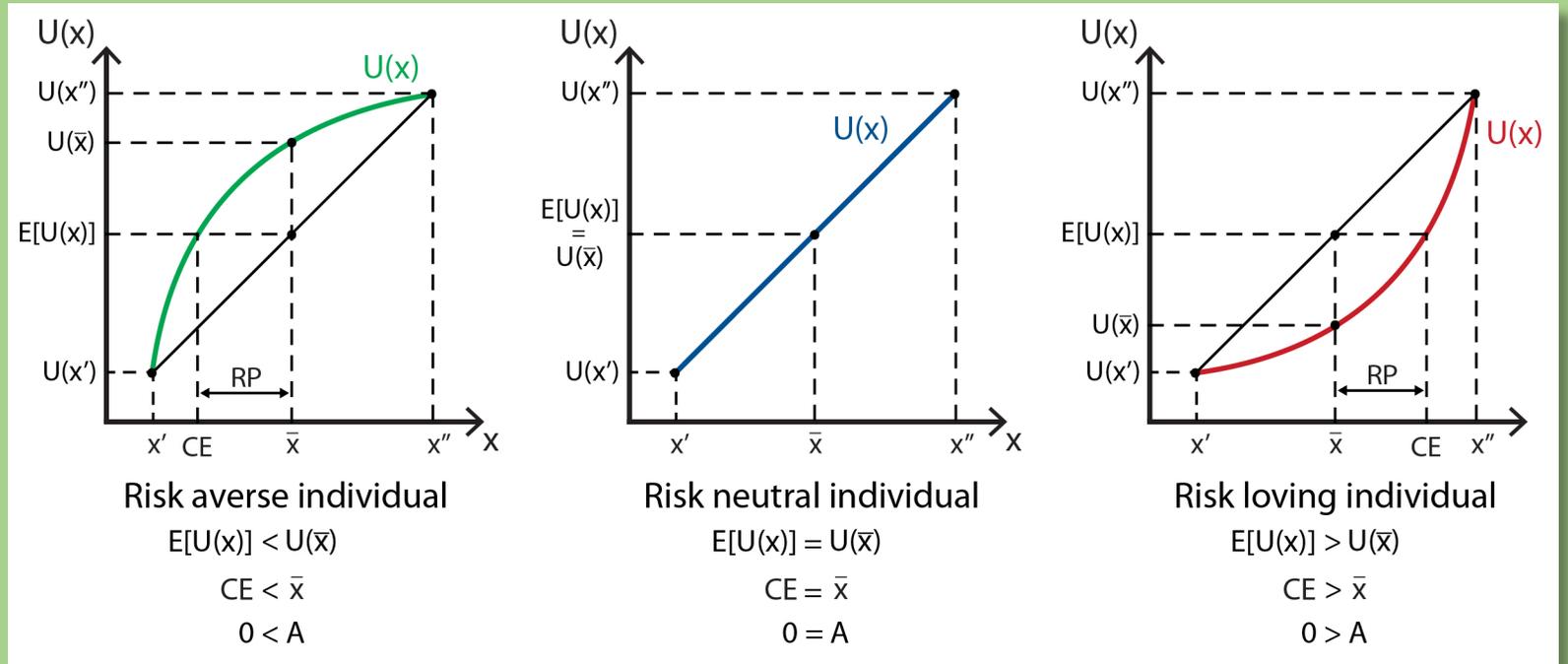
- In single-agent decision theory the key notion is that of an *optimal strategy*.
- In **noncooperative game theory**, the basic modeling unit is the **individual** (including its beliefs, preferences, and possible actions)
 - in **coalitional game theory**, it is the **group**
- Agents are **self-interested**
 - Each agent has its own description of which states of the world it **likes**—and that it **acts** in an attempt to **bring about** these states of the world.
- The dominant approach to modeling an agent's interests is **utility theory**
- The idea of **(expected) utility** can be grounded in a more basic concept of preferences **(and rationality): von Neumann and Morgenstern**

Why is “self-interestedness” crucial?

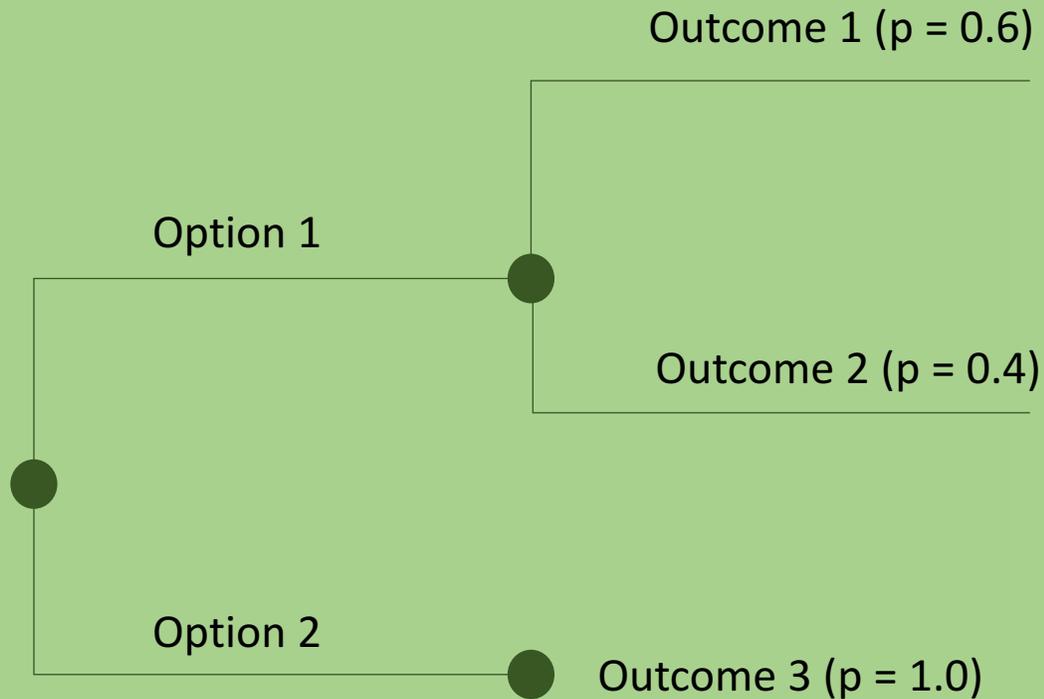


Utility

- How valuable is \$10 to you?
- \$100?
- \$10,000?
- \$100,000?
- \$101,000?
- \$102,000?
- ...



Expected Utility



Outcome 1 = \$1,000
Outcome 2 = \$100
Outcome 3 = \$640

Which option would you choose?



- What are $U_{Name}(\$1,000)$, $U_{Name}(\$640)$, $U_{Name}(\$100)$?
- What's the expected utility of Option 1? Option 2?

$$EU_{Name}(\text{Option1}) = 0.6 * U_{Name}(\$1,000) + 0.4 * U_{Name}(\$100)$$

$$EU_{Name}(\text{Option2}) = 1.0 * U_{Name}(\$640)$$

Rationality, risk attitudes



Utility Theory

- Let O denote a finite set of outcomes. For any pair $o_1, o_2 \in O$,
 - $o_1 \succcurlyeq o_2$: the agent weakly prefers o_1 to o_2
 - $o_1 \sim o_2$: the agent is indifferent between o_1 and o_2
 - $o_1 \succ o_2$: the agent strictly prefers o_1 to o_2 .
- Now, a **lottery** is the random selection of one of a set of outcomes according to specified probabilities.
- Formally, a lottery is a **probability distribution over outcomes** written $[p_1: o_1, \dots, p_k: o_k]$, where each $o_i \in O$, each $p_i \geq 0$ and $\sum_{i=1}^k p_i = 1$.
- Let L denote the set of all lotteries
 - Extend the \succcurlyeq relation to apply to the elements of L as well as to the elements of O , **effectively considering lotteries over outcomes to be outcomes themselves**

Utility Theory | Axioms

- **Axiom 3.1.1. Completeness.** $\forall o_1, o_2 \ o_1 \succ o_2 \text{ or } o_2 \succ o_1 \text{ or } o_1 \sim o_2.$
- **Axiom 3.1.2. Transitivity.** *If $o_1 \succcurlyeq o_2$ and $o_2 \succcurlyeq o_3$, then $o_1 \succcurlyeq o_3$.*
- **Axiom 3.1.3. Substitutability.** *If $o_1 \sim o_2$, then for all sequences of one or more outcomes o_3, \dots, o_k and sets of probabilities p, p_3, \dots, p_k for which $p + \sum_{i=3}^k p_i = 1$,*

$$[p : o_1, p_3 : o_3, \dots, p_k : o_k] \sim [p : o_2, p_3 : o_3, \dots, p_k : o_k].$$



Completeness: no ambiguity, preference must be there, can be partially ordered

Transitivity: consistency, inference

Substitutability states that **if an agent is indifferent between two outcomes, it is also indifferent between two lotteries that differ only in which of these outcomes is offered.**

Utility Theory | Axioms 2

- Let $P_\ell(o_i)$ denote the probability that outcome o_i is selected by lottery ℓ
- **Axiom 3.1.4 Decomposability.** *If $\forall o_i \in O, P_{\ell_1}(o_i) = P_{\ell_2}(o_i)$, then $\ell_1 \sim \ell_2$.*
 - **Decomposability states that an agent is always indifferent between lotteries that induce the same probabilities over outcomes**, no matter whether these probabilities are expressed through a single lottery or nested in a lottery over lotteries
- **Axiom 3.1.5 Monotonicity.** *If $o_1 \succ o_2$ and $p > q$, then $[p: o_1, 1 - p: o_2] \succ [q: o_1, 1 - q: o_2]$.*

Decomposability: Easier to compare lotteries
Monotonicity: consistent, dependable



Utility Theory | Axioms 3

- **Lemma 3.1.6** *If a preference relation \succsim satisfies the axioms **completeness**, **transitivity**, **decomposability**, and **monotonicity**, and if $o_1 \succ o_2$ and $o_2 \succ o_3$, then there exists some probability p such that*
 - for all $p' < p$, $o_2 \succ [p': o_1, (1 - p') : o_3]$, and*
 - for all $p'' > p$, $[p'': o_1, (1 - p'') : o_3] \succ o_2$*
- **Axiom 3.1.7 Continuity.** *If $o_1 \succ o_2$ and $o_2 \succ o_3$, then $\exists p \in [0, 1]$ such that $o_2 \sim [p: o_1, 1 - p: o_3]$.*

Utility Theory | Axioms 4

- If we accept Axioms 3.1.1, 3.1.2, 3.1.4, 3.1.5, and 3.1.7, it turns out that we have no choice but to accept the **existence of single-dimensional utility functions** whose expected values agents want to maximize
 - And if we do *not* want to reach this conclusion, we must then give up at least one of the axioms
- **Theorem 3.1.8 (von Neumann and Morgenstern, 1944)** *If a preference relation \succsim satisfies the axioms completeness, transitivity, substitutability, decomposability, monotonicity, and continuity, then there exists a function $u: O \rightarrow [0, 1]$ with the properties that*
 1. $u(o_1) \geq u(o_2)$ iff $o_1 \succsim o_2$, and
 2. $u([p_1: o_1, \dots, p_k: o_k]) = \sum_{i=1}^k p_i u(o_i)$

Normal Form Games

- **a.k.a. Strategic Form Games**
- **Definition 3.2.1 (Normal-form game)** A (finite, n -person) normal-form game is a tuple (N, A, u) , where:
 - N is a finite set of n players, indexed by i
 - $A = A_1 \times \cdots \times A_n$, where A_i is a finite set of actions available to player i
 - Each vector $a = (a_1, \dots, a_n) \in A$ is called an action profile
 - $u = (u_1, \dots, u_n)$ where $u_i: A_i \rightarrow \mathbb{R}$ is a real-valued utility (or payoff) function for player i
- Note that we previously argued that utility functions should map from the set of *outcomes*, not the set of *actions*
 - Here we make the implicit assumption that $O = A$.
- A natural way to represent games is via an n -dimensional matrix.

How should an agent play a normal form game?

- We have seen that under reasonable assumptions about preferences, agents will *always* have utility function whose expected values they want to maximize
- *This suggests that acting optimally in an uncertain environment is conceptually straightforward:*
- **Agents simply need to choose the course of action that maximizes expected utility!**
- *But in real life, that's often too good to be true*



Normal Form Games | Prisoner's Dilemma

	Player 2 No Betray	Player 2 Betray
Player 1 No Betray	1,1	-4,3
Player 1 Betray	3,-4	-3,-3

Best joint outcomes

Worst joint outcomes



Dominant strategy, Nash equilibrium

- Any rational user, when presented with this scenario **once**, will adopt Betray—regardless of what the other user does
 - Allowing the users to communicate beforehand will not change the outcome
- Perfectly rational agents will also adopt Betray even if they play **multiple** times
 - However, if the number of times that is infinite, or uncertain, agents may adopt No Betray

Normal Form Games | Common Payoff Game

- **Definition 3.2.2 (Common-payoff game)** A common-payoff game is a game in which for all action profiles $a \in A_1 \times \dots \times A_n$ and any pair of agents i, j , it is the case that $u_i(a) = u_j(a)$.
- a.k.a. **pure coordination games** or **team games**
 - no conflicting interests among agents; only need to coordinate to maximize benefits

E.g., two people walking towards each other in a hallway: If they choose the same side (left or right) they have some high utility; otherwise, a low utility.

	Agent 2 Left	Agent 2 Right
Agent 1 Left	1,1	0,0
Agent 1 Right	0,0	1,1



Why considering common payoff game? Agents need to understand on their own that they need to coordinate

Normal Form Games | Constant-Sum Game

- **Definition 3.2.3 (Constant-sum game)** A two-player normal-form game is constant sum if there exists a constant c such that for each strategy profile $a \in A_1 \times A_2$, it is the case that $u_1(a) + u_2(a) = c$.
- E.g., zero-sum game when $c = 0$
- Opposite of pure coordination games

	P2 Heads	P2 Tails
P1 Heads	1,-1	-1,1
P2 Tails	-1,1	1,-1



Paper-scissors-rock, constrained resource allocation (or assignment)

E.g., *Matching Pennies*. In this game, each of the two players has a penny and independently chooses to display either heads or tails. P1 pockets both if the pennies match; otherwise P2 pockets both

Connection to MAS?

Utility, expected utility, rationality?



Action selection, maximize utility, single agent decision making vs. multiple agents making decisions: **may not be able to secure "optimal" strategy**



Stupid Question: What if parking meters in a city allows for members of "car-pooling groups" to use left-over amount of money from parking meters, that would expire at the end of the day? What would your strategy be if you were a member of such a group?

