

# Introduction to the Boolean Satisfiability Problem

Spring 2020

CSCE 235H Introduction to Discrete Structures

URL: [cse.unl.edu/~cse235h](http://cse.unl.edu/~cse235h)

All questions: [Piazza](#)

# Satisfiability Study

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- 7 weeks
- 30 min lectures in recitation
- ~2 hours of homework per week
- Goals:
  - Exposure to fundamental research in CS
  - Understand how to model problems
  - Learn to use SAT solver, MiniSAT

# Boolean Satisfiability Problem

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- Given:
  - A Boolean formula
- Question:
  - Is there an assignment of truth values to the Boolean variables such that the formula holds true?

# Boolean Satisfiability Problem

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$$a \vee (\neg a \wedge b)$$

$$(a \vee \neg a) \rightarrow (b \wedge \neg b)$$

# Boolean Satisfiability Problem

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$$a \vee (\neg a \wedge b)$$

SATISFIABLE  
a=true, b=true

$$(a \vee \neg a) \rightarrow (b \wedge \neg b)$$

# Boolean Satisfiability Problem

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$$a \vee (\neg a \wedge b)$$

SATISFIABLE  
a=true, b=true

$$(a \vee \neg a) \rightarrow (b \wedge \neg b)$$

UNSATISFIABLE  
Left side of implication is a tautology.  
Right side of implication is a contradiction.  
True cannot imply false.

# Applications of SAT

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- Scheduling
- Resource allocation
- Hardware/software verification
- Planning
- Cryptography

# Conjunctive Normal Form

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- Variable  $a, b, p, q, x_1, x_2$
- Literal  $a, \neg a, q, \neg q, x_1, \neg x_1$
- Clause  $(a \vee \neg b \vee c)$
- Formula  $(a \vee \neg b \vee c) \wedge (b \vee c) \wedge (\neg a \vee \neg c)$



# Converting to CNF

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- All Boolean formulas can be converted to CNF
- The  $\rightarrow$ ,  $\leftrightarrow$ ,  $\oplus$  operators can be rewritten in terms of  $\neg$ ,  $\vee$ ,  $\wedge$
- $\neg$ ,  $\vee$ ,  $\wedge$  can be rearranged using
  - De Morgan's Laws
  - Distributive Laws
  - Double Negative
- May result in exponential size increase of the formula

# Converting to CNF

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$$(a \vee \neg a) \rightarrow (b \wedge \neg b) \equiv$$

# Converting to CNF

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Implication  $\neg(a \vee \neg a) \vee (b \wedge \neg b) \equiv$

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Implication  $\neg(a \vee \neg a) \vee (b \wedge \neg b) \equiv$

DeMorgan's  $(\neg a \wedge a) \vee (b \wedge \neg b) \equiv$

# Converting to CNF

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$$(a \vee \neg a) \rightarrow (b \wedge \neg b) \equiv$$

Implication  $\neg(a \vee \neg a) \vee (b \wedge \neg b) \equiv$

DeMorgan's  $(\neg a \wedge a) \vee (b \wedge \neg b) \equiv$

$$(\neg a \vee b) \wedge (\neg a \vee \neg b) \wedge (a \vee b) \wedge (a \vee \neg b)$$

Distributive

# Interpretation of CNF

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- Every clause must be satisfied by at least one true literal
- Total possible number of solutions increases as number of variables increases
- Clauses constrain the possible solutions
- Smaller clauses are more constraining

# Interpretation of CNF

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$$(a \vee \neg b \vee \neg c) \\ \wedge (b \vee c) \\ \wedge (\neg a)$$

a	b	c
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1

# Interpretation of CNF

---

$$\begin{aligned} (a \vee \neg b \vee \neg c) &\equiv \\ \wedge (b \vee c) &\rightarrow \neg(\neg a \wedge b \wedge c) \\ \wedge (\neg a) & \end{aligned}$$

a	b	c
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1



# Interpretation of CNF

$$(a \vee \neg b \vee \neg c)$$

$$\wedge (b \vee c) \equiv \neg(\neg b \wedge \neg c)$$

$$\wedge (\neg a)$$

a	b	c
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1

# Interpretation of CNF

$$(a \vee \neg b \vee \neg c)$$

$$\wedge (b \vee c)$$

$$\wedge (\neg a)$$

$\equiv$

$$\neg(a)$$

a	b	c
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1

# Interpretation of CNF

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$$(a \vee \neg b \vee \neg c) \\ \wedge (b \vee c) \\ \wedge (\neg a)$$

a	b	c
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1	0	0
1	0	1
1	1	0
1	1	1

# Determining SAT/UNSAT

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- All that is required to show satisfiability is to find a valid solution
- Many techniques available:
  - Guessing and checking
  - Systematic search
  - Inference

# Systematic Search with Backtracking

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- Construct a binary tree of all combinations
- Proceeds in a depth first manner
- Each level corresponds to a variable
- Each branch corresponds to a truth assignment
- Branches of the tree are 'pruned' when the assignment cannot be extended in a satisfiable manner

# Systematic Search with Backtracking

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$$(a \vee b \vee c)$$
$$\wedge(\neg a \vee \neg b)$$
$$\wedge(\neg b \vee \neg c)$$
$$\wedge(\neg c \vee \neg a)$$

root

# Systematic Search with Backtracking

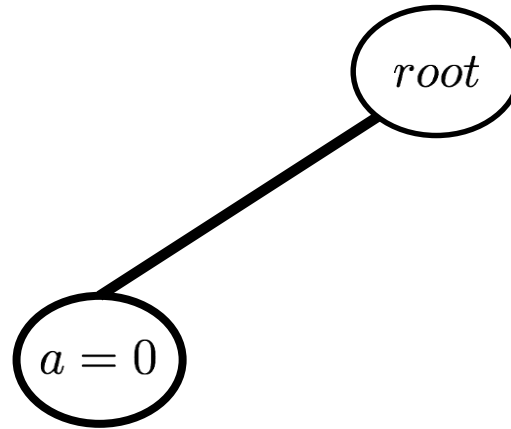
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$$(\underline{a} \vee b \vee c)$$

$$\wedge(\underline{\neg a} \vee \neg b)$$

$$\wedge(\neg b \vee \neg c)$$

$$\wedge(\neg c \vee \underline{\neg a})$$



# Systematic Search with Backtracking

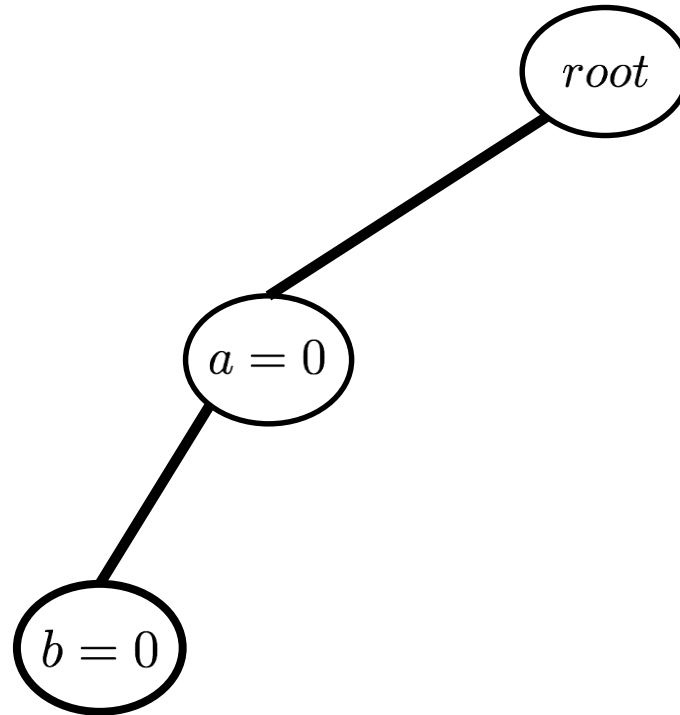
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$$(\underline{a} \vee \underline{b} \vee c)$$

$$\wedge(\underline{\neg a} \vee \underline{\neg b})$$

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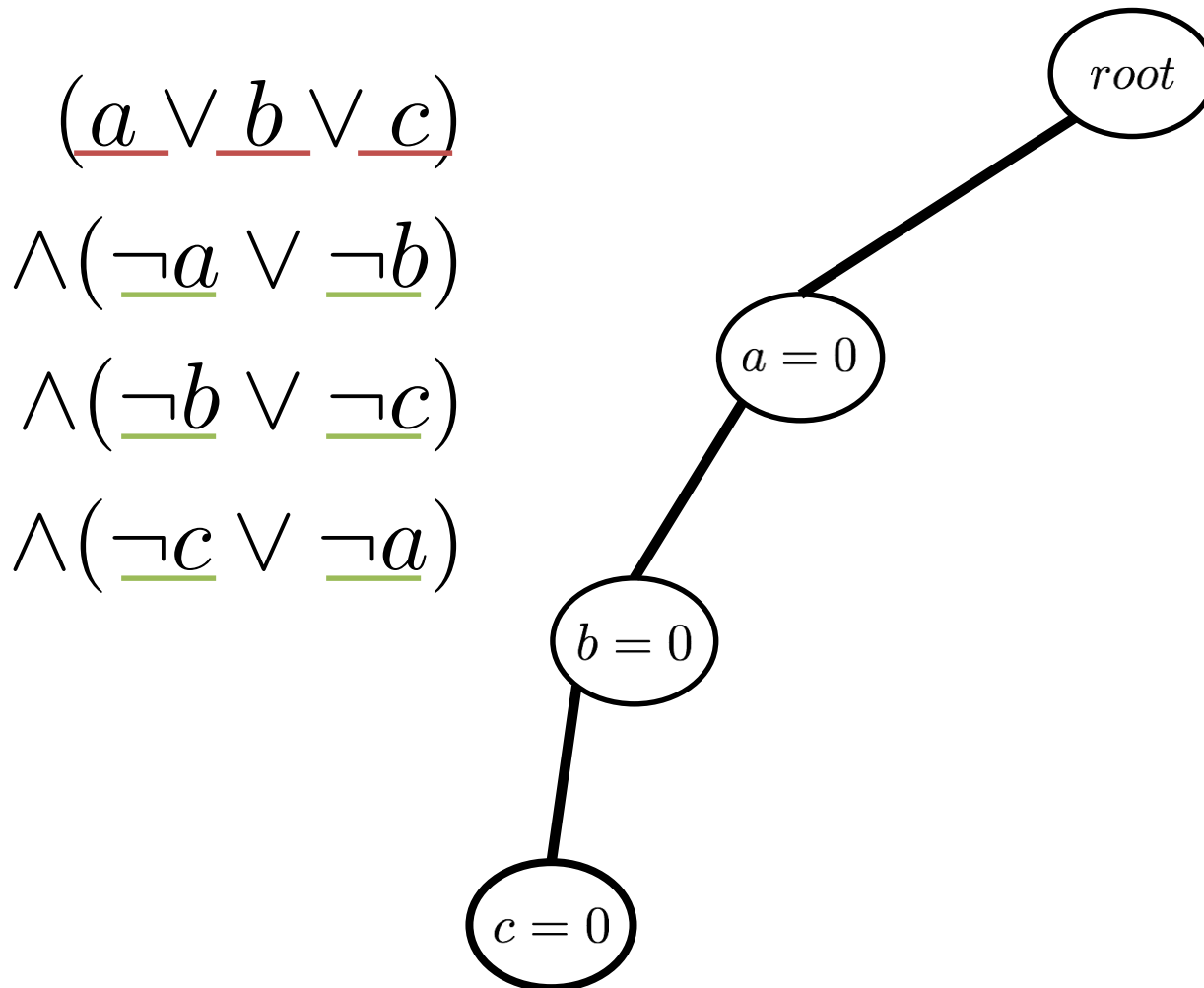
$$\wedge(\neg c \vee \underline{\neg a})$$





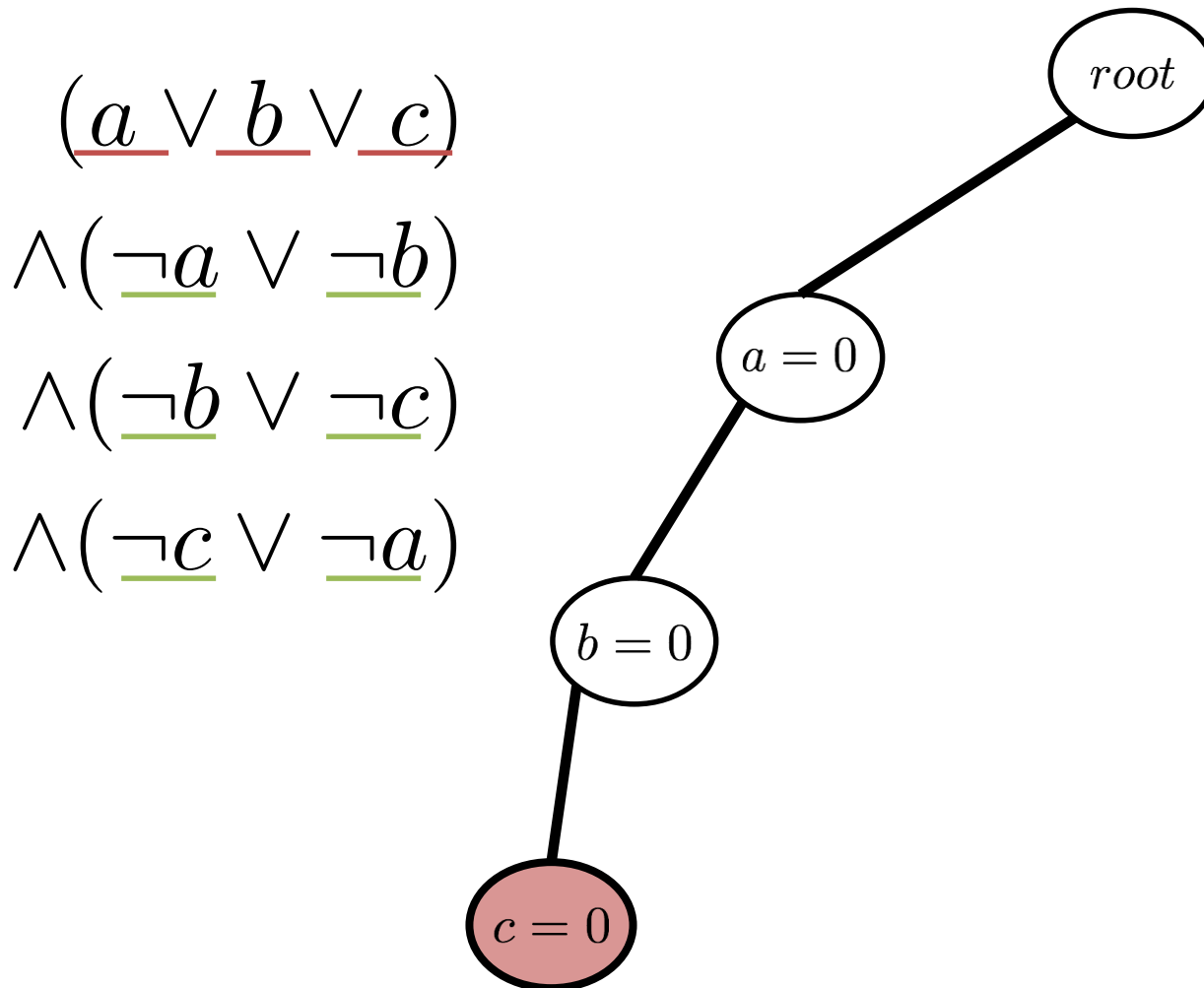
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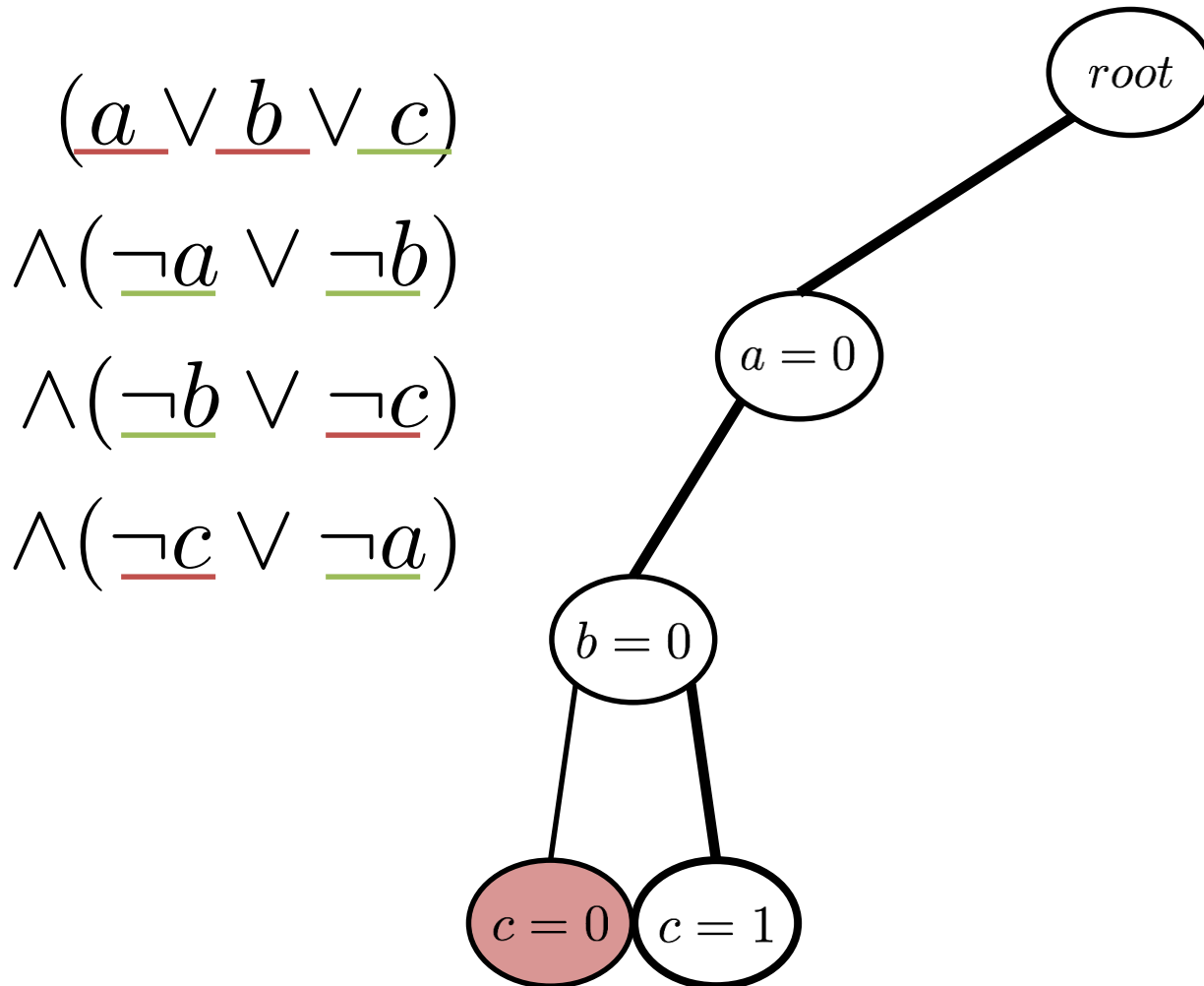
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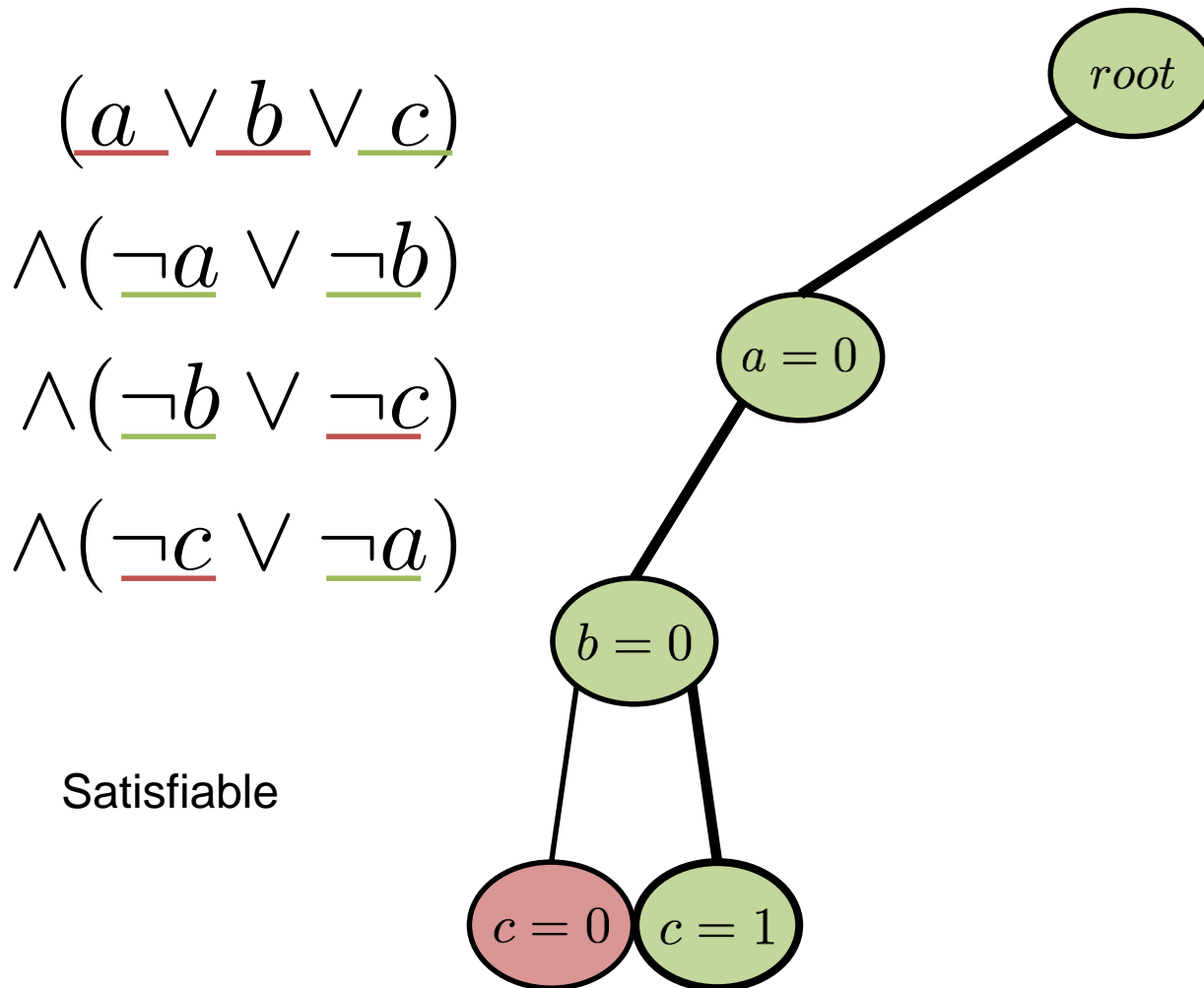


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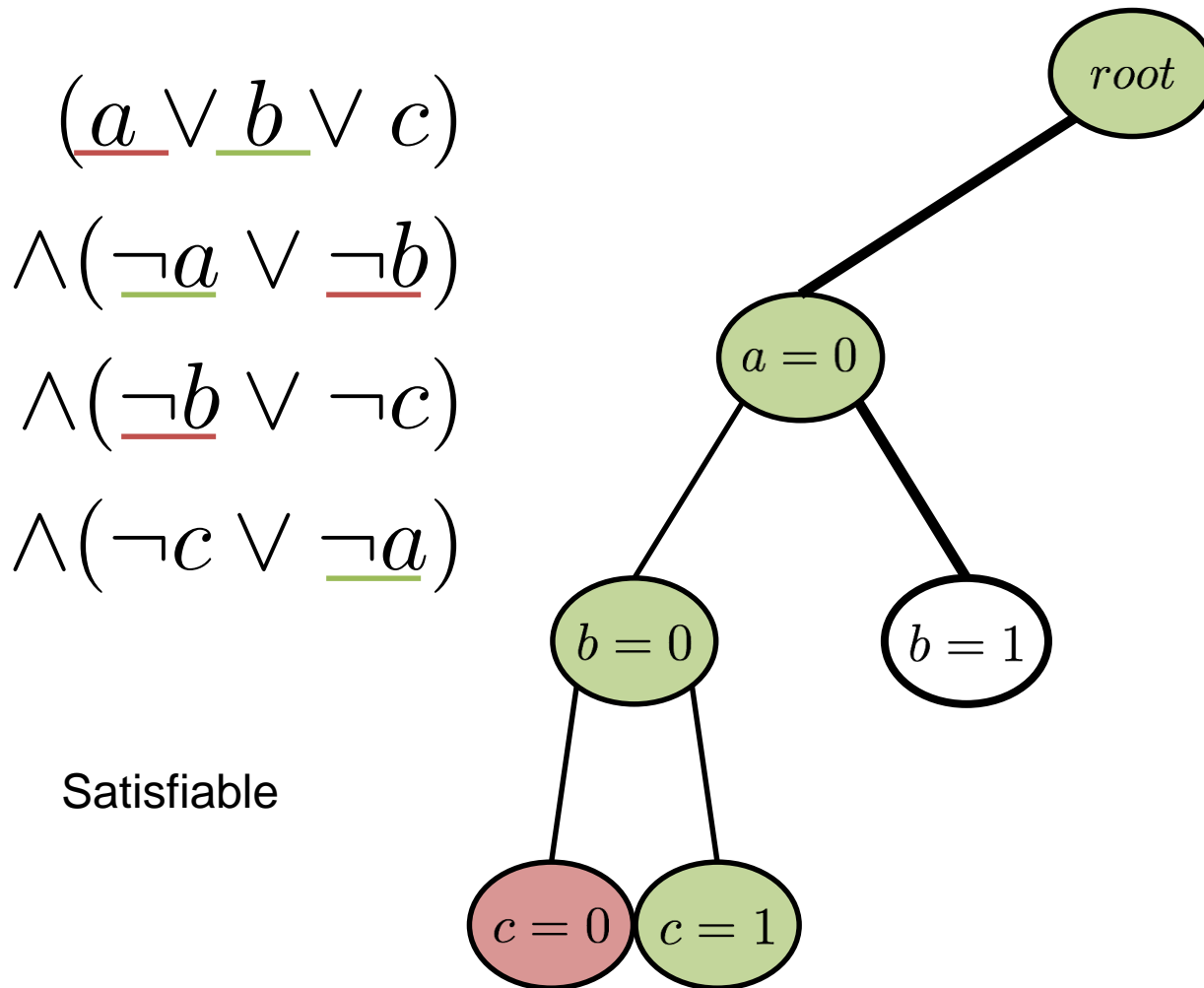
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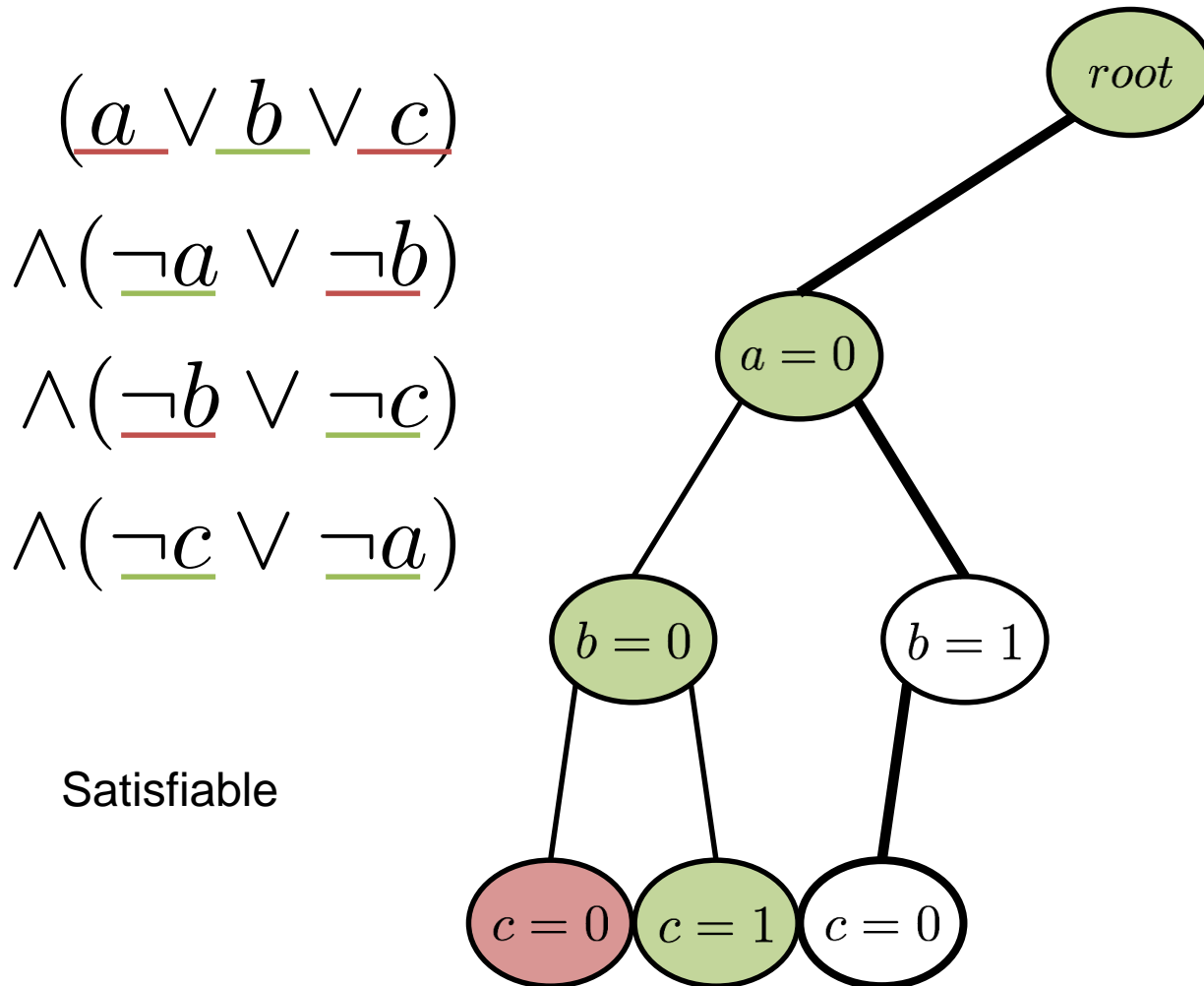
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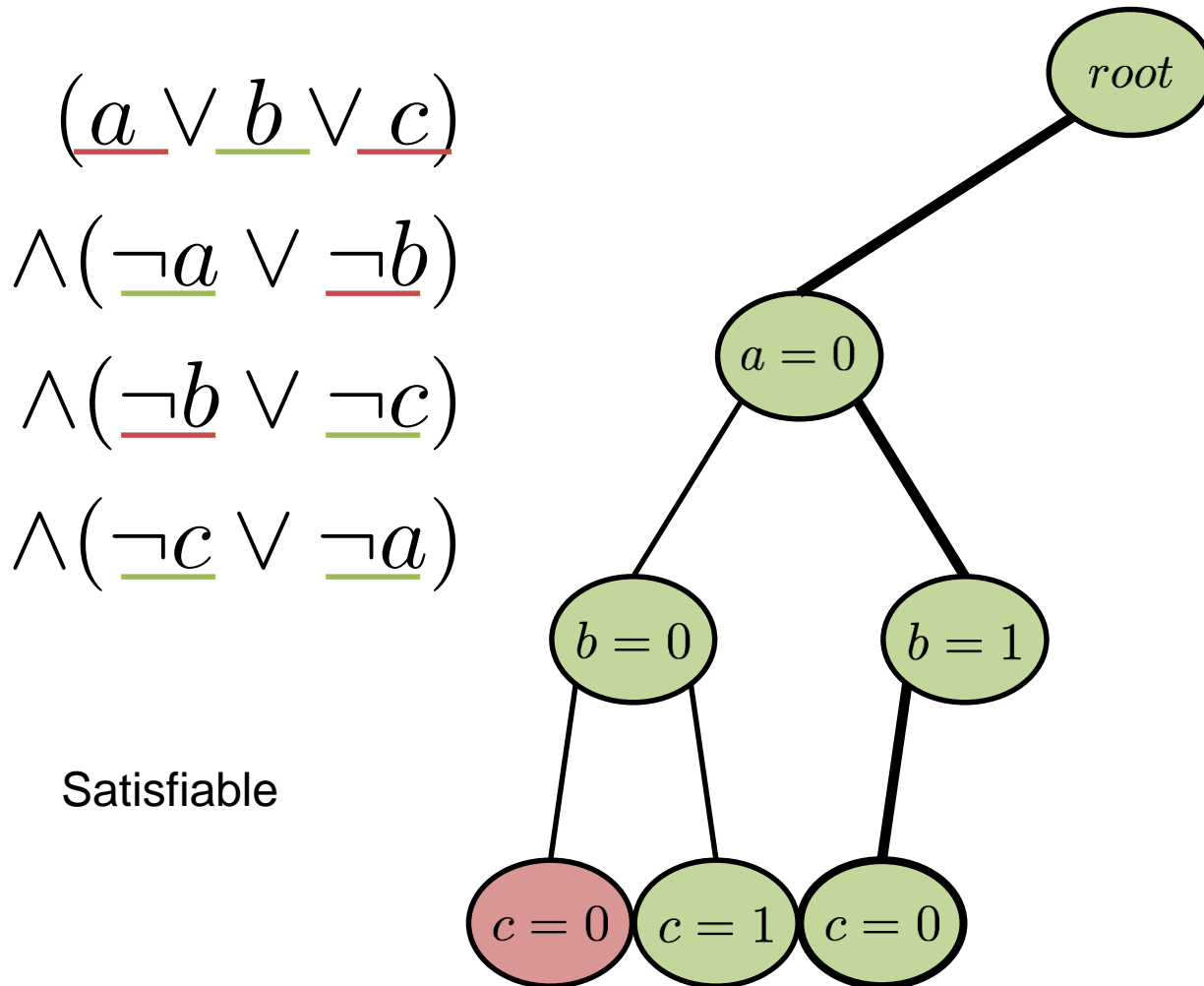
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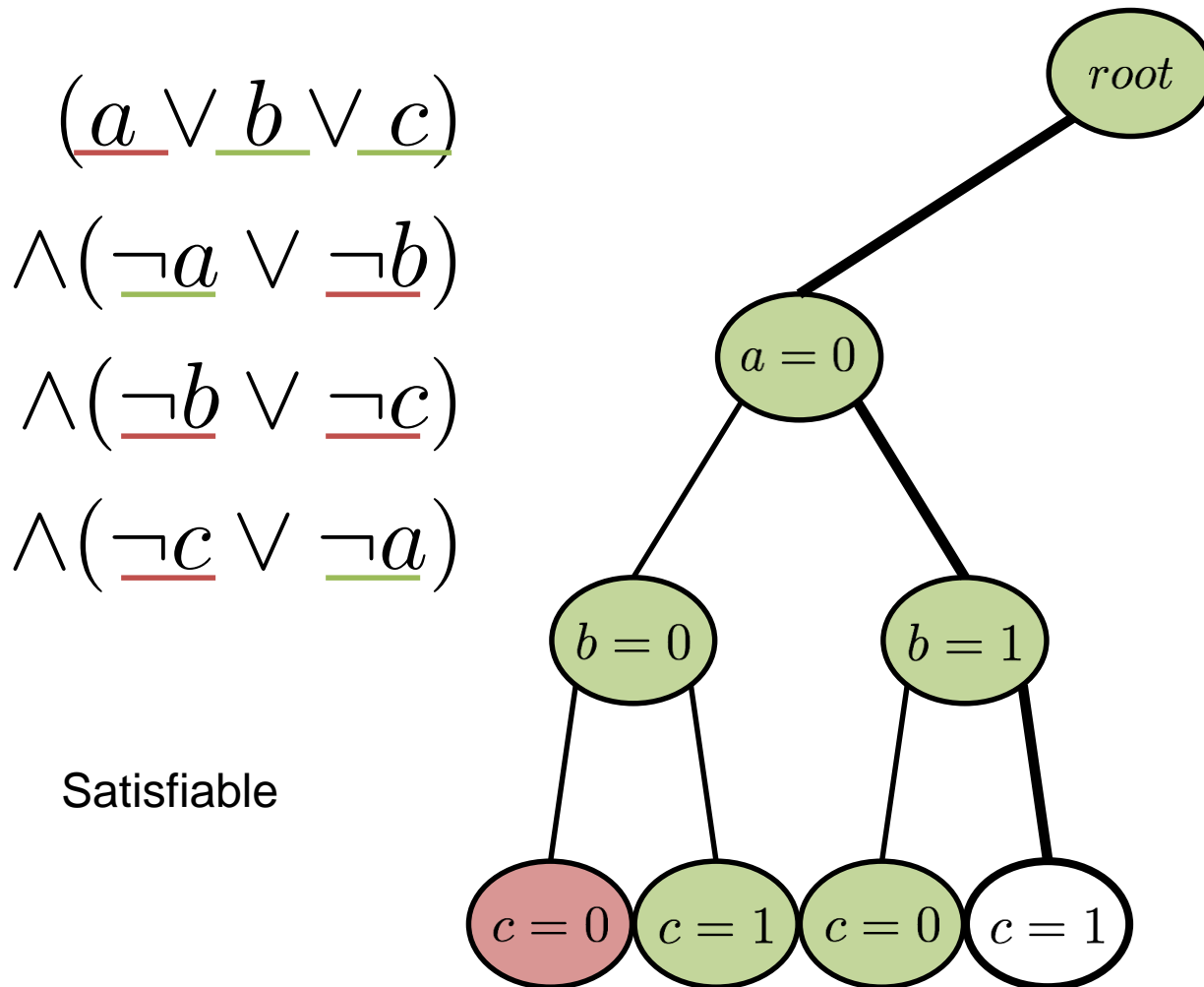
# Systematic Search with Backtracking



# Systematic Search with Backtracking

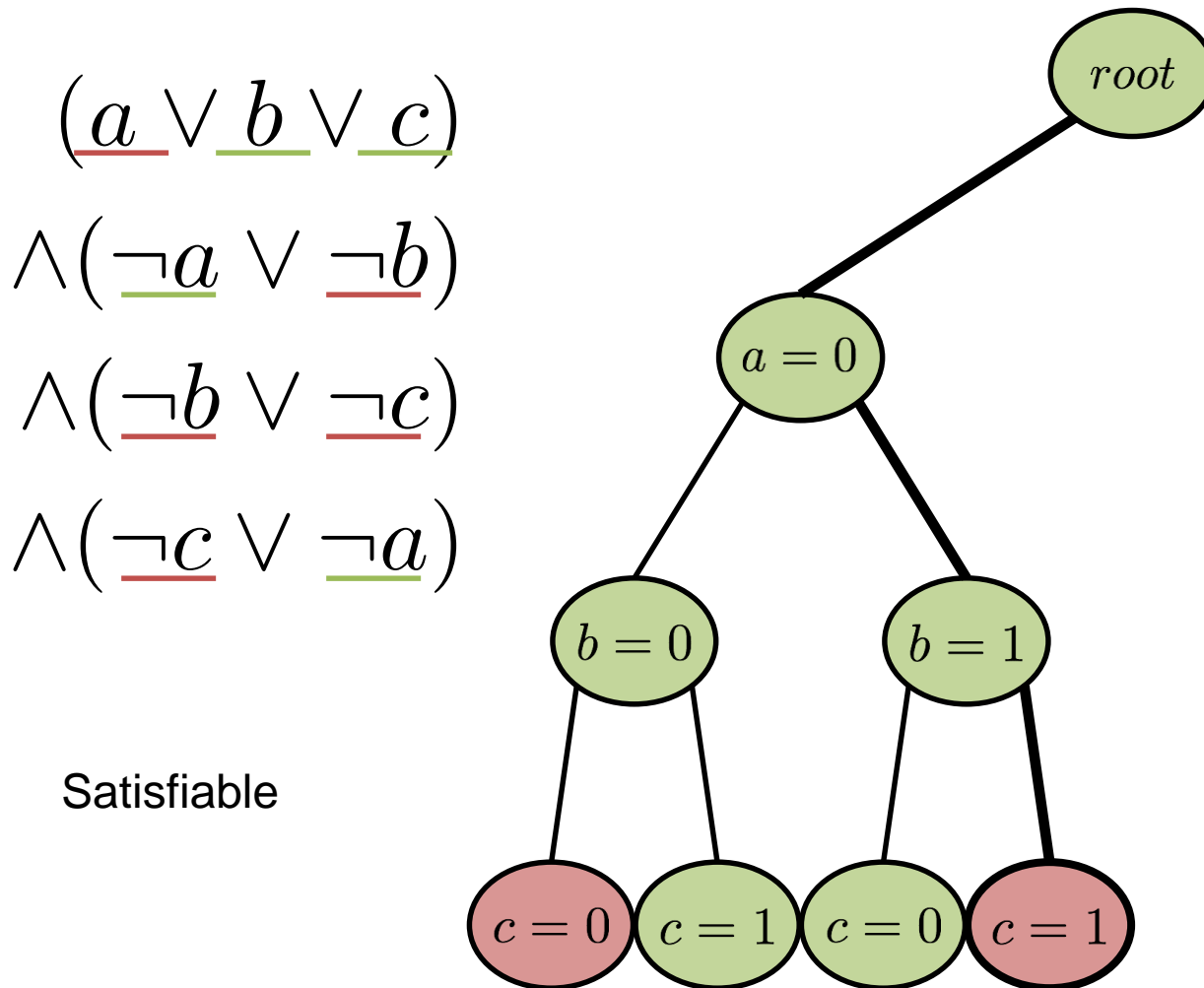


# Systematic Search with Backtracking





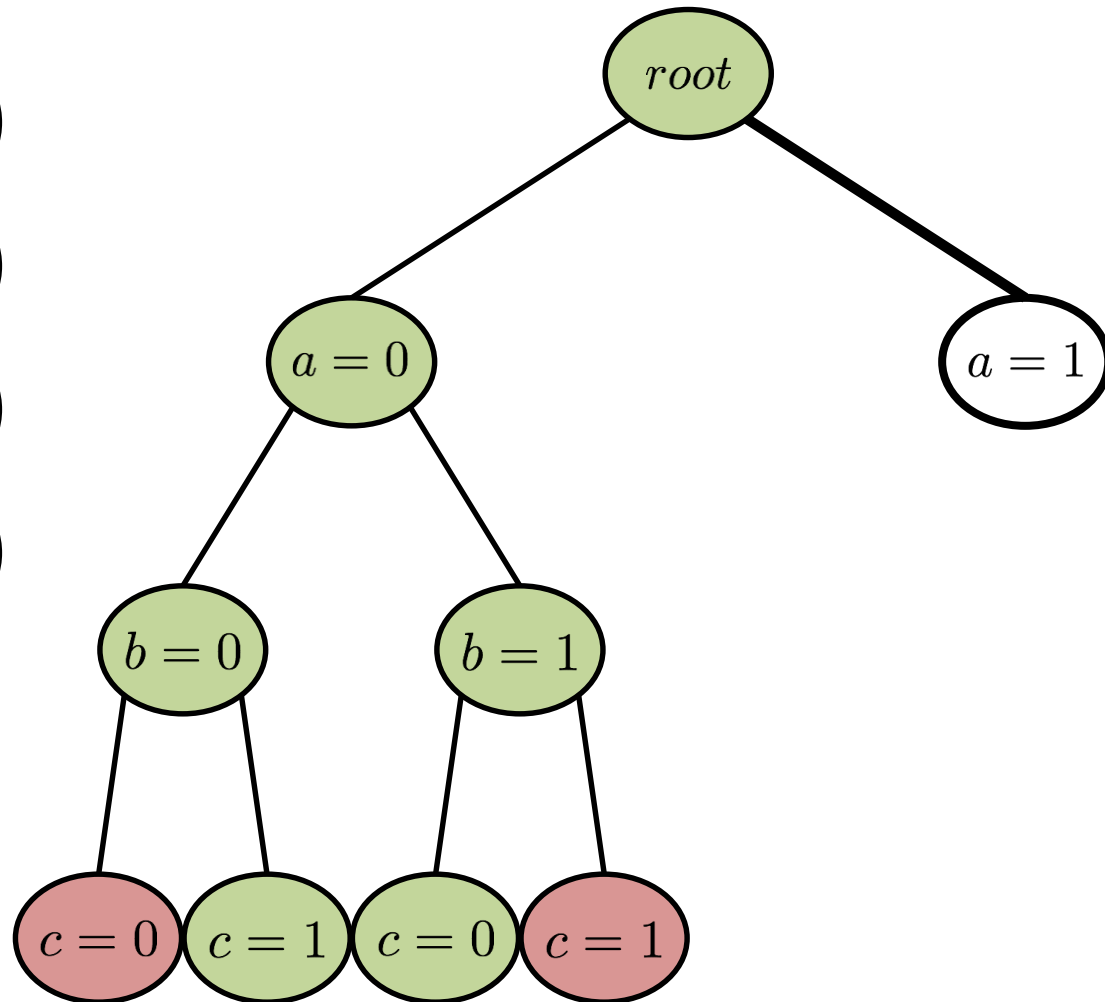
# Systematic Search with Backtracking



# Systematic Search with Backtracking

$$\begin{aligned} &(\underline{a} \vee b \vee c) \\ &\wedge(\underline{\neg a} \vee \neg b) \\ &\wedge(\neg b \vee \neg c) \\ &\wedge(\neg c \vee \underline{\neg a}) \end{aligned}$$

Satisfiable



# Systematic Search with Backtracking

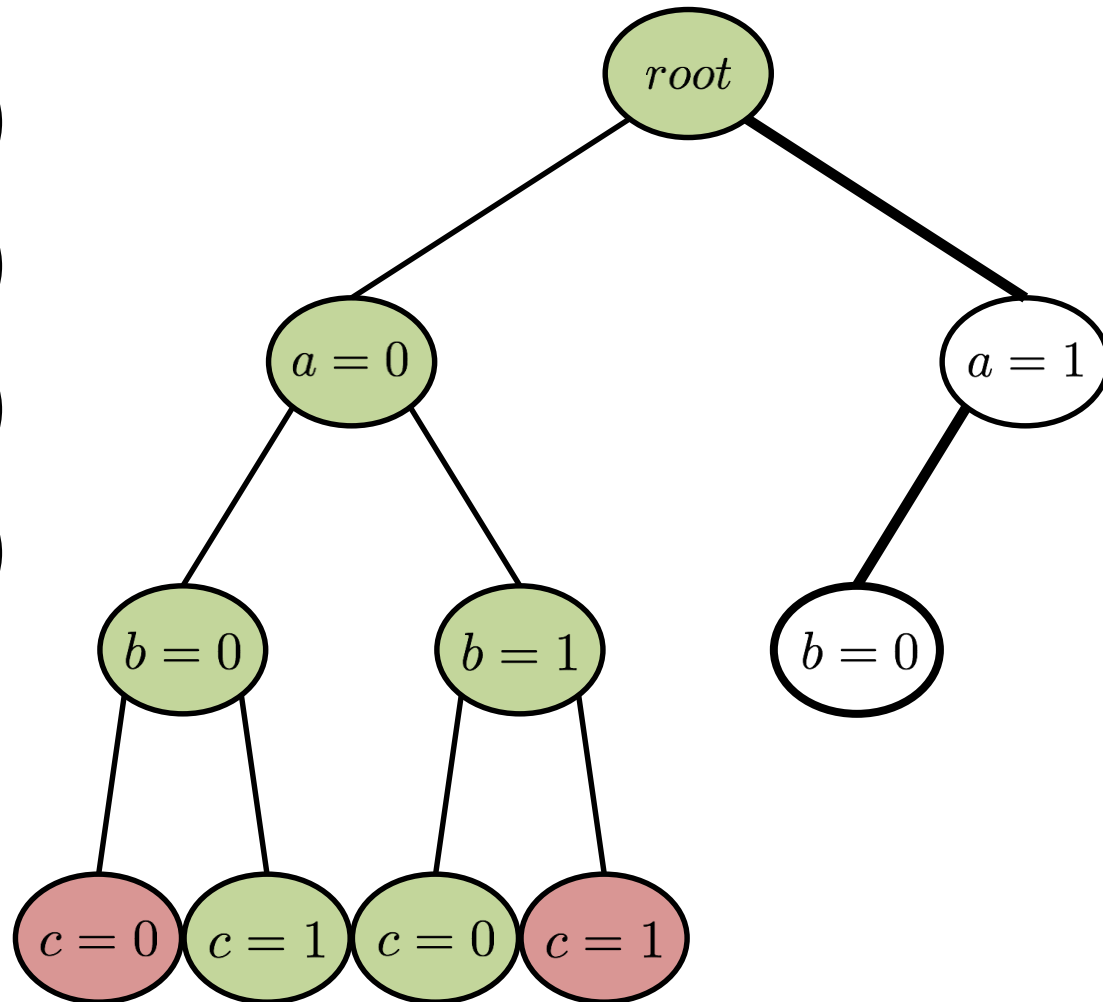
$$(\underline{a} \vee \underline{b} \vee c)$$

$$\wedge(\underline{\neg a} \vee \underline{\neg b})$$

$$\wedge(\underline{\neg b} \vee \neg c)$$

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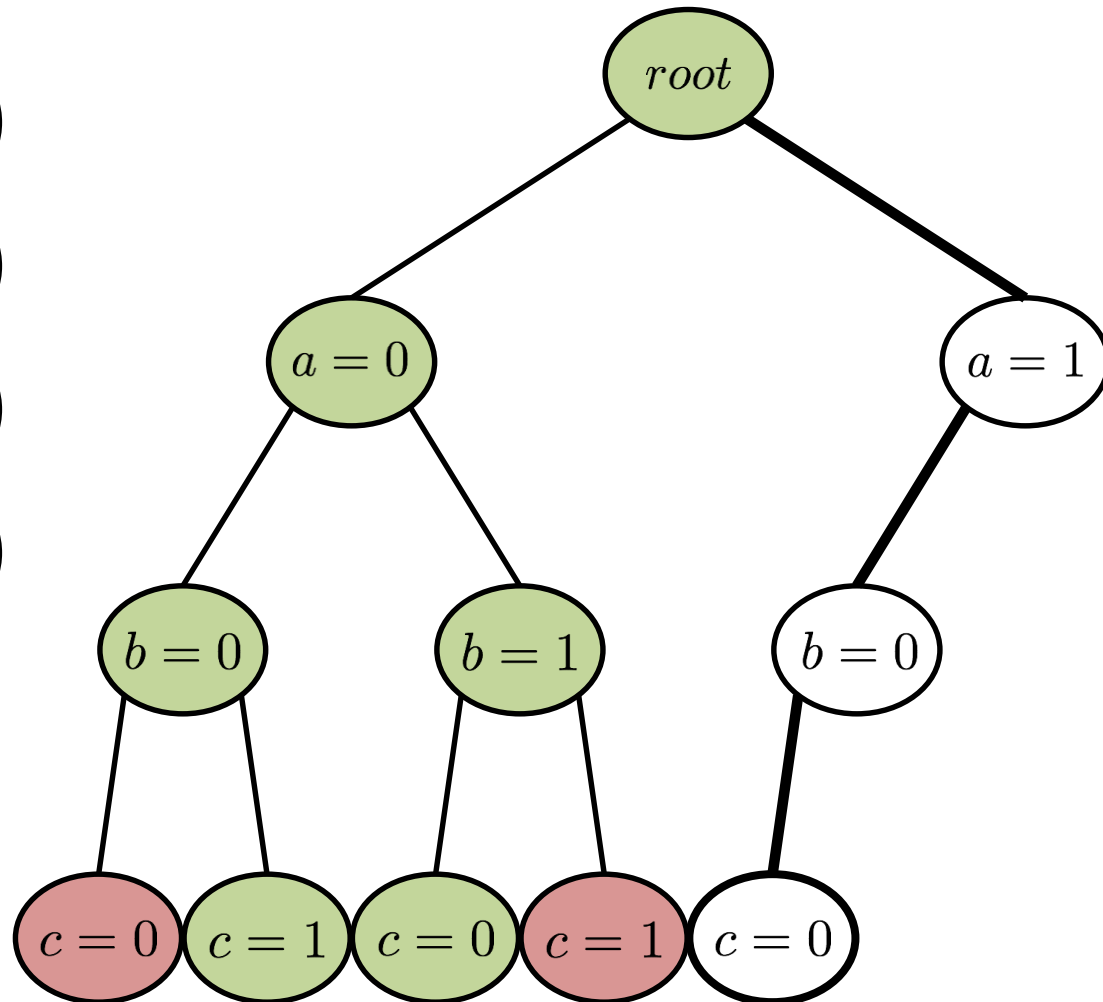
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# Systematic Search with Backtracking

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Satisfiable



# Systematic Search with Backtracking

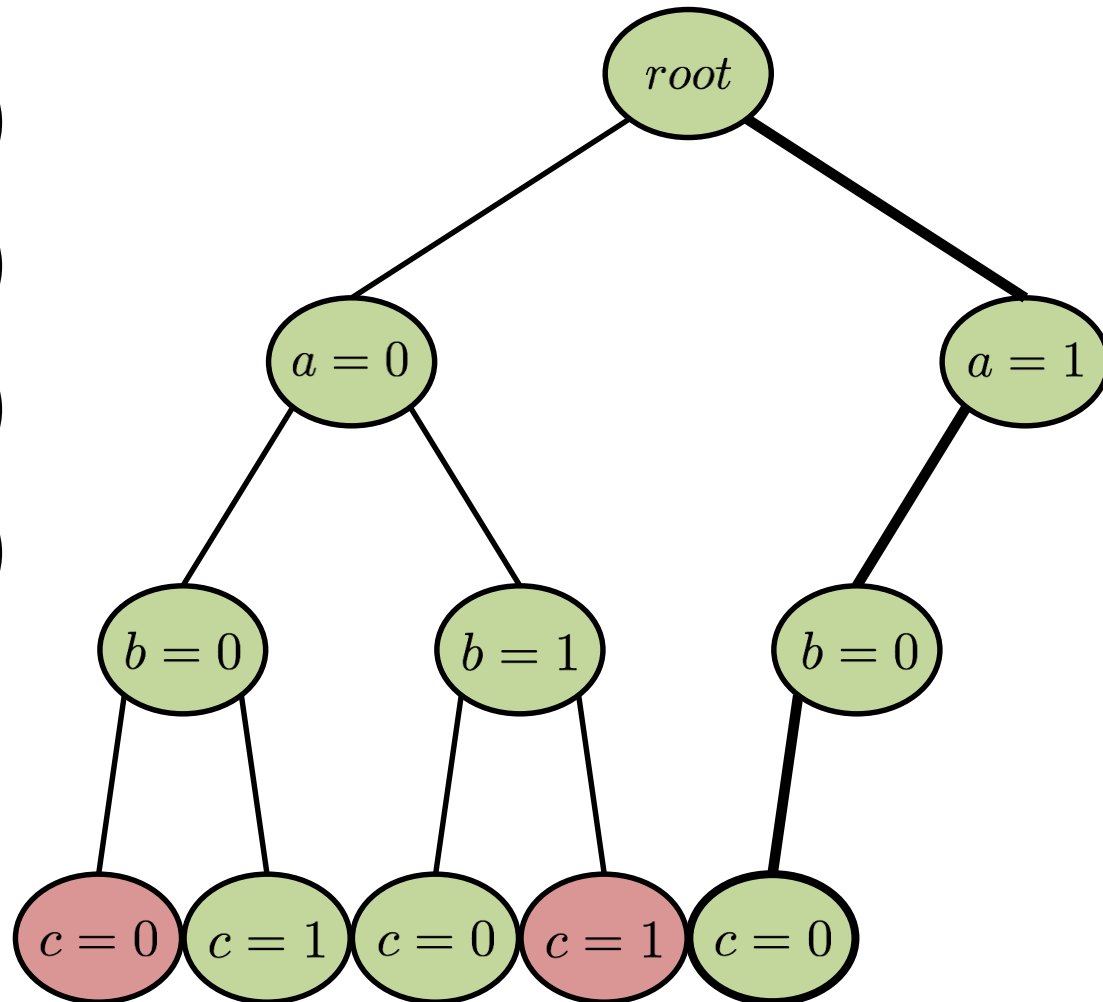
$$(\underline{a} \vee \underline{b} \vee \underline{c})$$

$$\wedge(\underline{\neg a} \vee \underline{\neg b})$$

$$\wedge(\underline{\neg b} \vee \underline{\neg c})$$

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Satisfiable



# Systematic Search with Backtracking

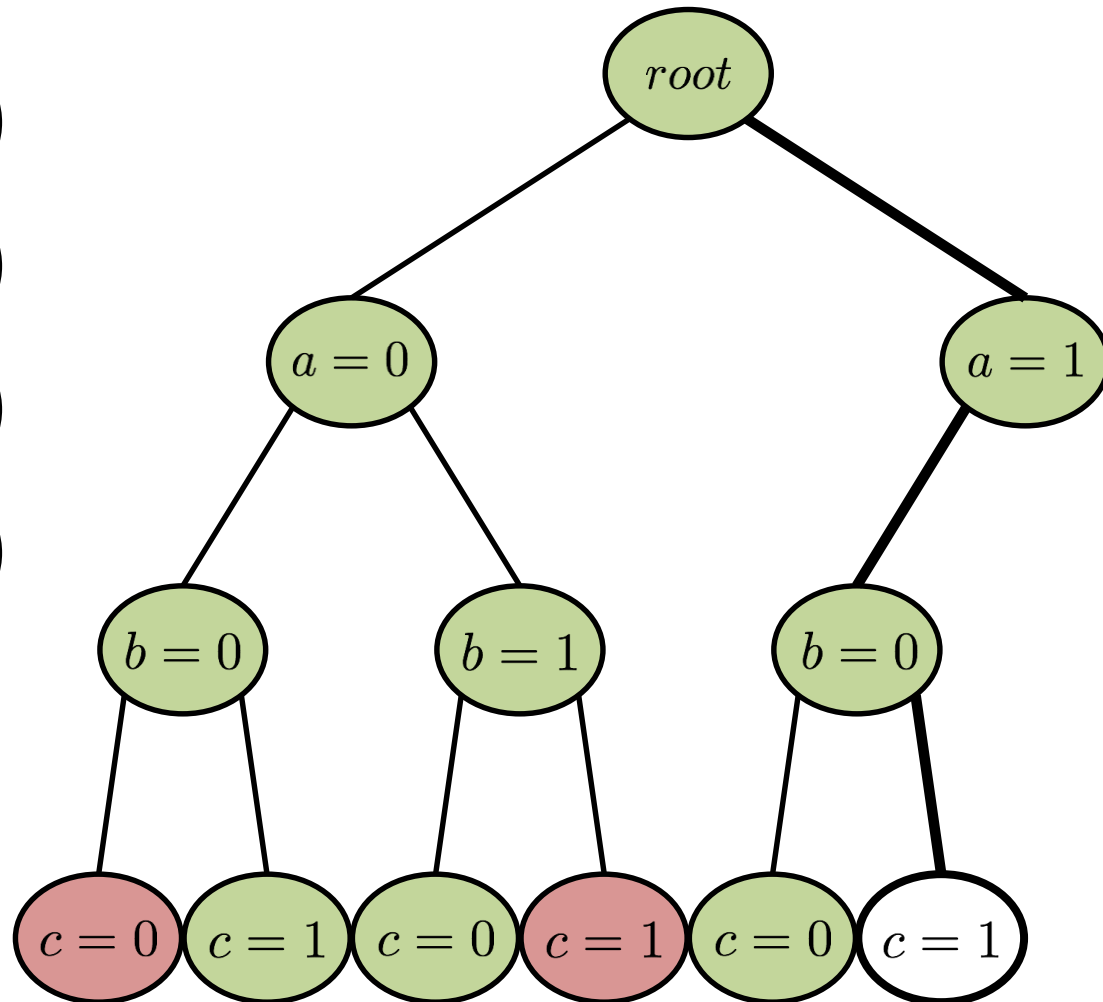
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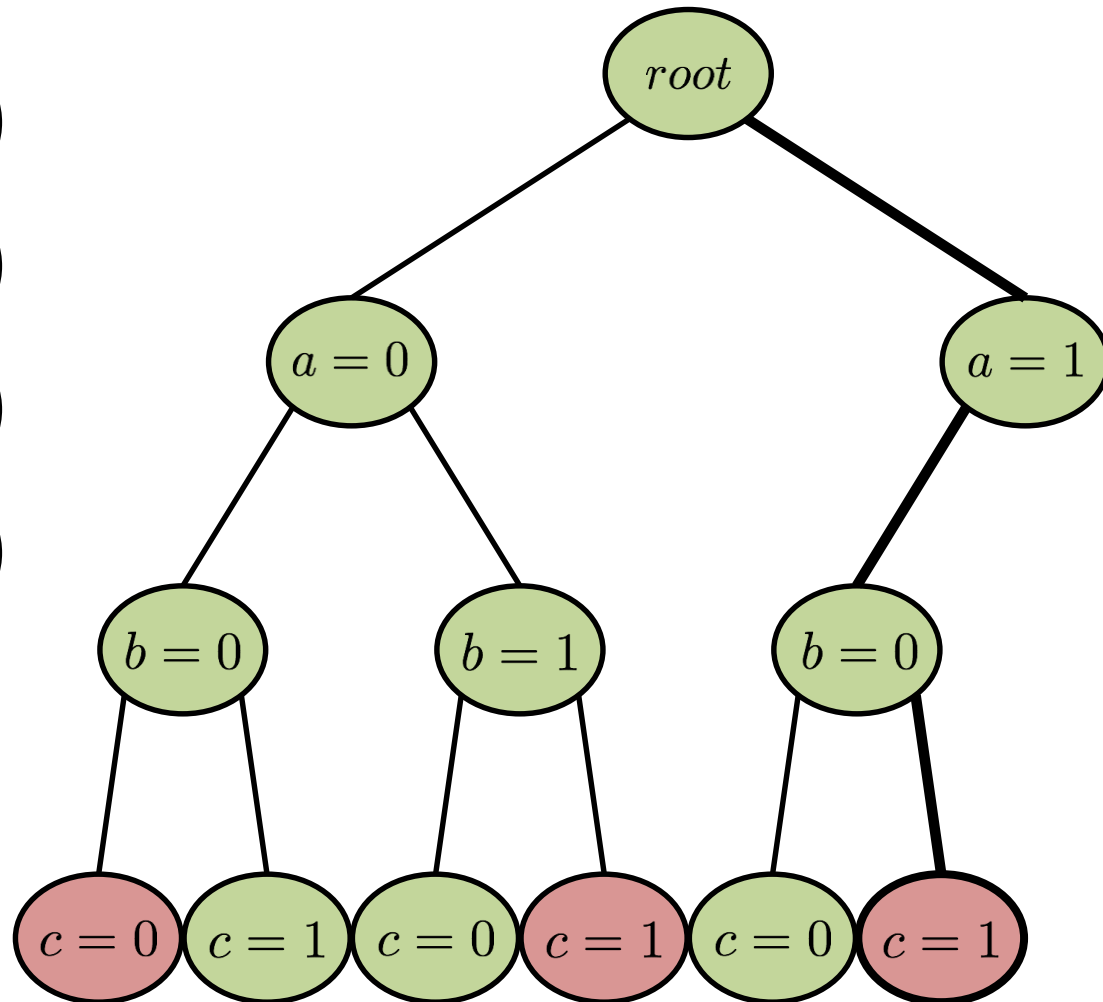
Satisfiable



# Systematic Search with Backtracking

$$\begin{aligned} &(\underline{a} \vee \underline{b} \vee \underline{c}) \\ &\wedge(\underline{\neg a} \vee \underline{\neg b}) \\ &\wedge(\underline{\neg b} \vee \underline{\neg c}) \\ &\wedge(\underline{\neg c} \vee \underline{\neg a}) \end{aligned}$$

Satisfiable



# Systematic Search with Backtracking

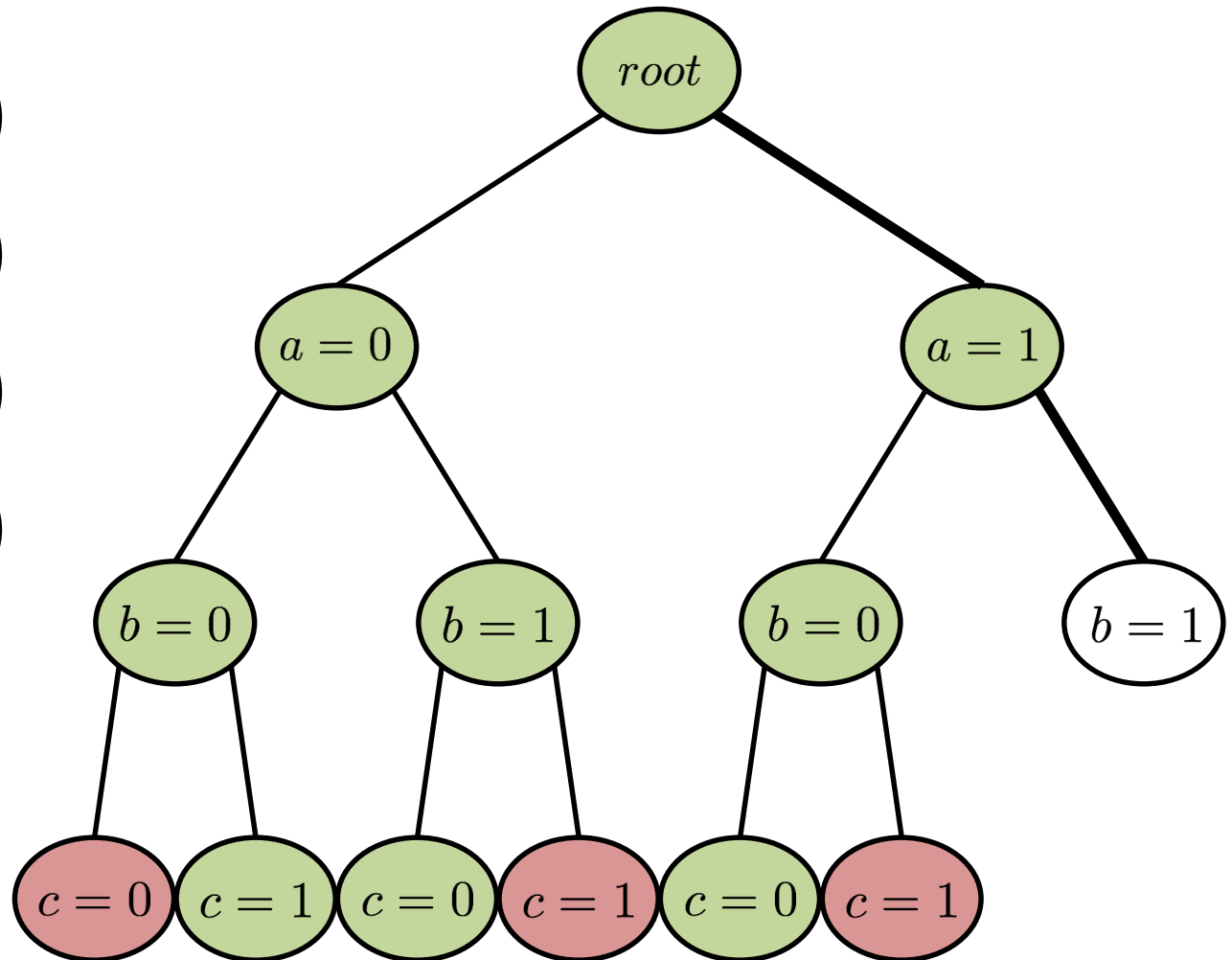
$$(\underline{a} \vee \underline{b} \vee c)$$

$$\wedge(\underline{\neg a} \vee \underline{\neg b})$$

$$\wedge(\underline{\neg b} \vee \neg c)$$

$$\wedge(\neg c \vee \underline{\neg a})$$

Satisfiable

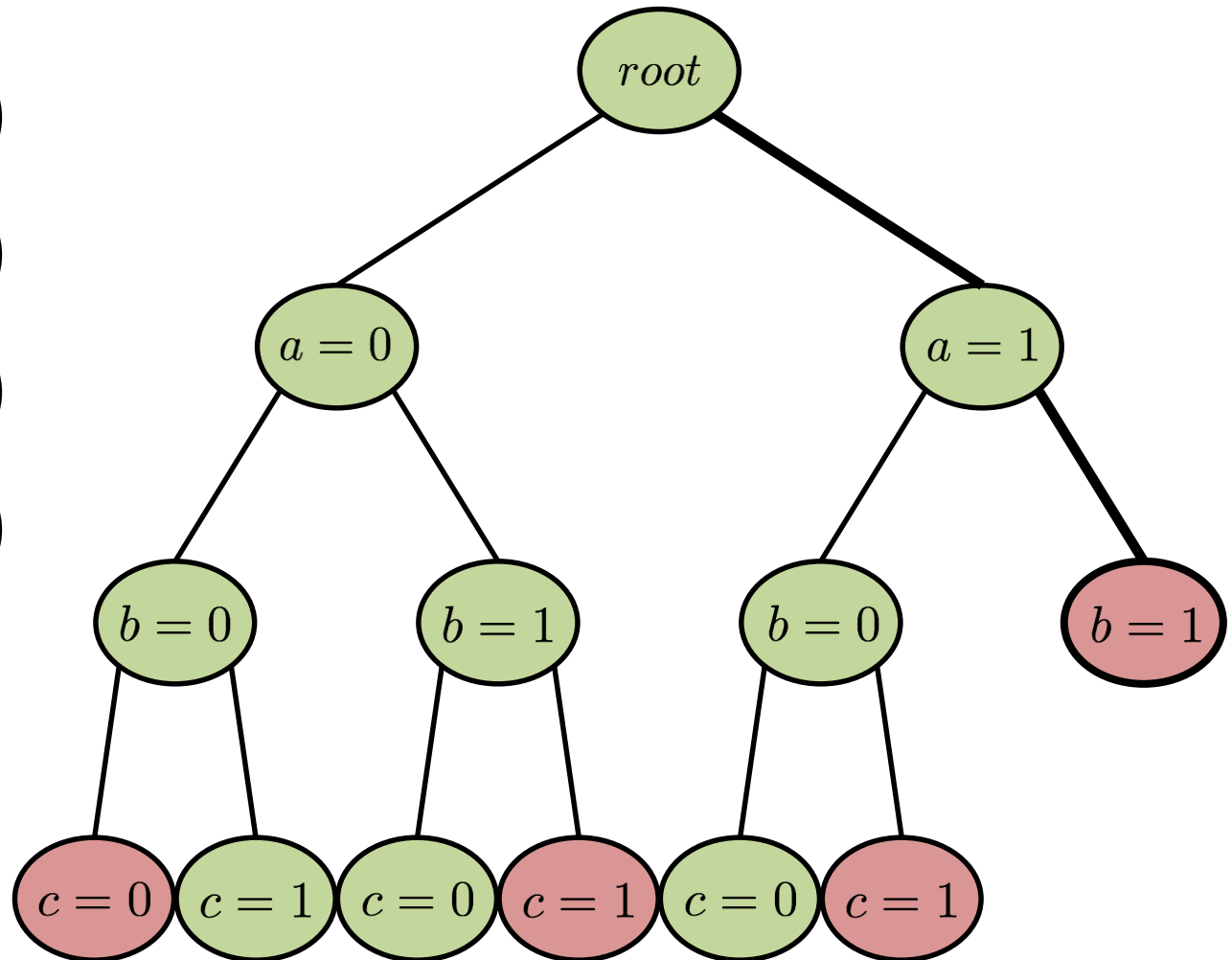




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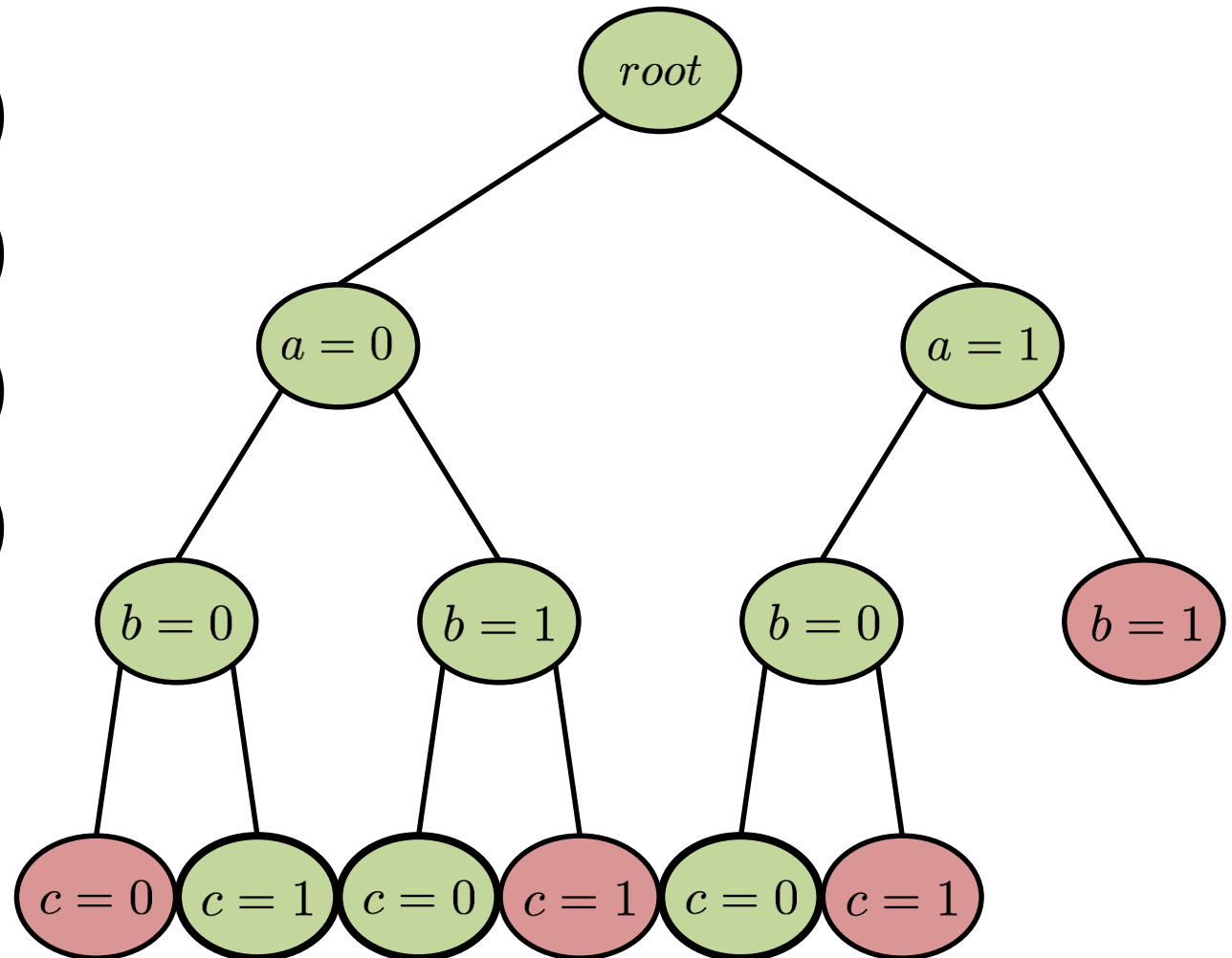
$$\begin{aligned} &(\underline{a} \vee \underline{b} \vee c) \\ &\wedge(\underline{\neg a} \vee \underline{\neg b}) \\ &\wedge(\underline{\neg b} \vee \neg c) \\ &\wedge(\neg c \vee \underline{\neg a}) \end{aligned}$$

Satisfiable



# Systematic Search with Backtracking

$$(a \vee b \vee c)$$
$$\wedge (\neg a \vee \neg b)$$
$$\wedge (\neg b \vee \neg c)$$
$$\wedge (\neg c \vee \neg a)$$



Satisfiable  
3 solutions total