

Unit Propagation and Variable Ordering in MiniSAT

CSCE 235H Introduction to Discrete Structures

Spring 2019

URL: cse.unl.edu/~cse235h

All questions: Piazza

Outline

- Unit propagation
- Unit resolution
- Thrashing calls for a good variable ordering
- Variable ordering by
 - Activity
 - Decay

Unit Propagation (1): In a clause

- A satisfied clause can be removed

- If $a \leftarrow true$, remove all clauses where it is positive (positive occurrences)

$$a \vee b \vee c$$

- If $a \leftarrow false$, remove all clauses where it is negative (negative occurrences)

$$\neg a \vee b \vee c$$

- A clause can be simplified

- If $a \leftarrow false$, remove it from all clauses where it is positive

$$a \vee b \vee c \qquad b \vee c$$

- If $a \leftarrow true$, remove all clauses where it is negative

$$\neg a \vee b \vee c \qquad b \vee c$$

Unit Propagation (2)

- When a literal in a clause is true, the entire clause is true and can be removed

After assignment: $a \leftarrow false$

$$\begin{aligned} & \dots \wedge C_i \wedge (\neg a \vee b \vee \neg c) \wedge C_{i+2} \wedge \dots \\ \equiv & \dots \wedge C_i \wedge (true \vee b \vee \neg c) \wedge C_{i+2} \wedge \dots \\ \equiv & \dots \wedge C_i \wedge true \wedge C_{i+2} \wedge \dots \\ \equiv & \dots \wedge C_i \wedge C_{i+2} \wedge \dots \end{aligned}$$

Unit Propagation (3)

- When a literal in a clause is false, the literal may be removed from the clause

After assignment: $a \leftarrow true$

$$(\neg a \vee b \vee \neg c)$$

$$\equiv (false \vee b \vee \neg c)$$

$$\equiv (b \vee \neg c)$$

Unit Propagation (4)

- Assignments may result in unit clauses
 $a \leftarrow false$ and $a \vee b$ yield b
- Unit clauses immediately force an assignment
 $b \leftarrow true$
- This can lead to a chain reaction as new assignments ‘propagate’ throughout the clauses

Unit Propagation: Example

- SAT formula with 1,000 variables

- Cycle of implications

$$v_1 \rightarrow v_2$$

$$v_2 \rightarrow v_3$$

$$\vdots$$

- A single assignment results in unit propagation to solve the entire problem

$$(\neg v_1 \vee v_2)$$

$$\wedge (\neg v_2 \vee v_3)$$

$$\vdots$$

$$\wedge (\neg v_{999} \vee v_{1000})$$

$$\wedge (\neg v_{1000} \vee v_1)$$

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Unit Resolution

- Unit resolution or disjunctive syllogism

$$\frac{a, \neg a \vee b}{b}$$

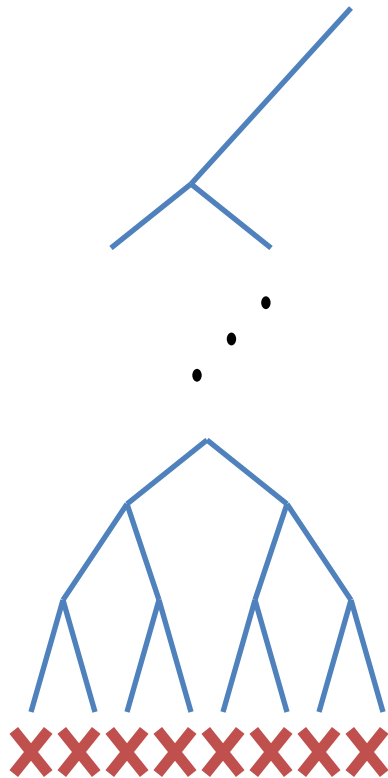
- Unit propagation has similarities to unit resolution
 - Assignments add unit clauses to the formula
 - When a new unit clause is added, it can be resolved with other clauses in the formula and the result added to the formula

$$(a) \wedge (\neg a \vee b \vee \neg c) \rightarrow (b \vee \neg c)$$

Outline

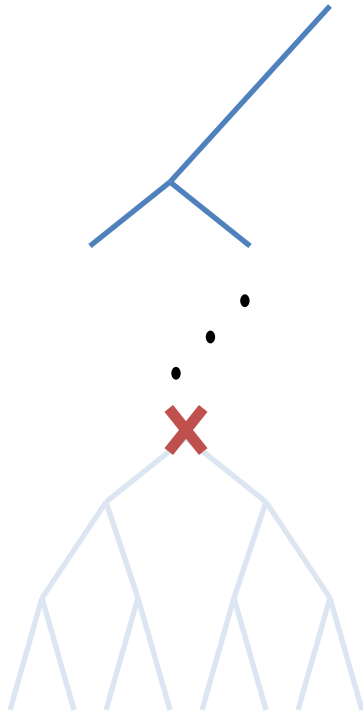
- Unit propagation
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Thrashing in Search



- Repeatedly hitting conflict within a subtree of the search tree
- Thrashing is costly

Fail-First Principle (FFP)



- It is better to fail early than to waste time exploring a subtree with no solutions
- Want to handle potential conflicts earlier

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Variable Ordering Heuristics

- The order in which variables are assigned
- Heuristic: good performance in general, not guaranteed to be optimal
- Static versus dynamic
 - Static: Assignment order is decided before search and maintained fixed throughout search
 - Dynamic: Variable ordering is adjusted during the course of the search

Activity-Based Heuristics

- Variables are assigned an ‘activity’ value
- Variables involved in a conflict have their activities increased
- Activities exponentially decay
- Select most ‘active’ variable

$$(\neg a \vee \neg c \vee d) \times \rightarrow$$

$$Activity(a)_+ = 1$$

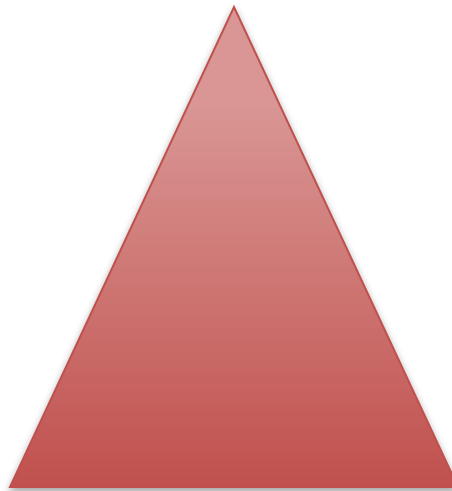
$$Activity(c)_+ = 1$$

$$Activity(d)_+ = 1$$

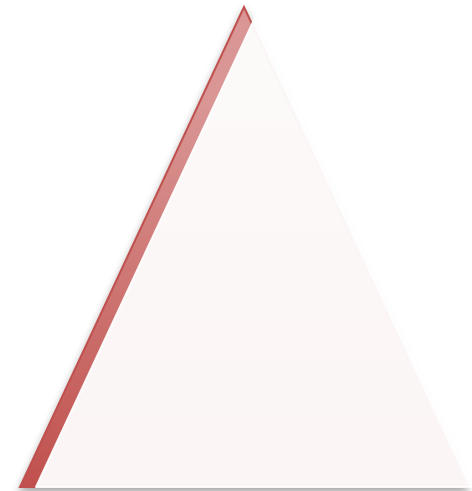
Activity-Based Heuristics Example

- UNSAT formula with 1000 variables
- Conflict occurs between v_{999} and v_{1000}

$$\begin{aligned} & \vdots \\ & \wedge (v_{999} \vee v_{1000}) \\ & \wedge (\neg v_{999} \vee v_{1000}) \\ & \wedge (v_{999} \vee \neg v_{1000}) \\ & \wedge (\neg v_{999} \vee \neg v_{1000}) \end{aligned}$$



Lexicographic ordering



Activity-based ordering