Recitation 11: Asymptotics and Summations

Created by Taylor Spangler, Adapted by Beau Christ

April 15, 2018

- Problem 3.2:25: Give a good big-O estimate for the following functions:
 - $(n^{2}+8)(n+1) = n^{3} + n^{2} + 8n + 1$. We can find n_{0}, c to prove that it is $O(n^{3})$
 - $(n \log n + n^2)(n^3 + 1) = n^5 + n^4 \log n + n^2 + n \log n$, again easily this is $O(n^5)$, perhaps using the limit method.
 - $(n! + 2^n)(n^3 + \log(n^2 + 1)) = n! \cdot n^3 + n! \cdot \log(n^2 + 1) + 2^n \cdot n^3 + 2^n \cdot \log(n^2 + 1).$ Well, this one is a little bit trickier. It is actually $O(n! \cdot n^3)$, why is it not O(n!) or 2^n ?
- Problem 3.2:31: Show that

$$f(x) \in \Theta(g(x)) \Leftrightarrow f(x) \in O(g(x)) \land g(x) \in O(f(x))$$

- 1. \Rightarrow First we begin with a definition: $f(x) \in \Theta(g(x))$ if $f(x) \in O(g(x))$ and $f(x) \in \Omega(g(x))$
 - So we can see that $\exists c_1, c_2 \in (R)^+$ such that $f(x) \leq c_1 \cdot g(x)$. We also know that $f(x) \geq c_2 \cdot g(x)$.
 - Reversing the second inequality we get $g(x) \leq c \cdot f(x)$ where $c = \frac{1}{c^2}$.
 - So then we have $f(x) \in O(g(x))$ and $g(x) \in O(f(x))$
- $2. \Leftarrow$
 - Suppose $f(x) \in O(g(x))$ and $g(x) \in O(f(x))$.
 - Then we know that $\exists c_1, c_2$ such that $f(x) \leq c_1 \cdot g(x)$ and $g(x) \leq c_2 \cdot f(x)$.
 - so $\frac{1}{c_1}g(x) \leq f(x) \leq c_2 \cdot g(x)$.
 - But this is the definition of Θ , therefore $f(x) \in \Theta(g(x))$.
- What is the tightest bound we can form here:
 - 1. $x^2 + 3x + 5 \in \Delta(x^3)$: big-O
 - 2. $2^n \log(6) + n^2 \in \Delta(2^n)$: Θ
 - 3. $2^n \cdot n! + 2^n \log(n) \in \Delta(2^n)$: Ω

• To prove the first one of the previous questions, we use the limit method We have:

$$\lim_{x \to \infty} \frac{x^2 + 3x + 4}{x^3} = \lim_{x \to \infty} \frac{2x + 3}{3x^2} = \lim_{x \to \infty} \frac{2}{6x} = 0$$

Therefore we can conclude that x^3 grows much faster, and so we get that $x^2 + 3x + 4 \in O(x^3)$

- Next up Sequences: Can we name the first 4 terms of the following sequence $\{2^n + 1\}_{n=0}^{\infty}$?
 - 1. $a_0 = 2^0 + 1 = 2$
 - 2. $a_1 = 2^1 + 1 = 3$ 3. $a_2 = 2^2 + 1 = 5$
 - 5. $u_2 2 + 1 = 0$
 - 4. $a_3 = 2^3 + 1 = 9$

• Now compute the following sum $\sum_{i=1}^{5} 6$ We have:

$$\sum_{i=1}^{5} 6 = 6 \sum_{i=1}^{5} 1 = 6 \cdot (5 - 1 + 1) = 6 \cdot 5 = 30$$

• How about the following geometric $\sum_{i=1}^{8} 3 \cdot 2^{i}$

- We can recognize that this expression is the sum of a geometric progression (i.e., geometric series), which is in general given as $\sum_{i=0}^{n} ar^{i}$, except that here we start at 1 instead of zero, so we can simply compute it by subtracting the first term from the sum.
- The formula for computing this is $\frac{a \cdot r^{n+1}-a}{r-1}$, here r=2 and a=3
- Using the formula we can get $\frac{3\cdot 2^9-3}{1} = 3\cdot 512$ 3 = 1533
- However, remember we have to subtract the first term, so $1533 3 \cdot 2^0 = 1530$.