

# Recitation 11: Asymptotics and Summations

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- Problem 3.2:25: Give a good big- $O$  estimate for the following functions:

- $(n^2 + 8)(n + 1) = n^3 + n^2 + 8n + 1$ . We can find  $n_0, c$  to prove that it is  $O(n^3)$
- $(n \log n + n^2)(n^3 + 1) = n^5 + n^4 \log n + n^2 + n \log n$ , again easily this is  $O(n^5)$ , perhaps using the limit method.
- $(n! + 2^n)(n^3 + \log(n^2 + 1)) = n! \cdot n^3 + n! \cdot \log(n^2 + 1) + 2^n \cdot n^3 + 2^n \cdot \log(n^2 + 1)$ . Well, this one is a little bit trickier. It is actually  $O(n! \cdot n^3)$ , why is it not  $O(n!)$  or  $2^n$ ?

- Problem 3.2:31: Show that

$$f(x) \in \Theta(g(x)) \Leftrightarrow f(x) \in O(g(x)) \wedge g(x) \in O(f(x))$$

1.  $\Rightarrow$  First we begin with a definition:  $f(x) \in \Theta(g(x))$  if  $f(x) \in O(g(x))$  and  $f(x) \in \Omega(g(x))$

- So we can see that  $\exists c_1, c_2 \in (R)^+$  such that  $f(x) \leq c_1 \cdot g(x)$ . We also know that  $f(x) \geq c_2 \cdot g(x)$ .
- Reversing the second inequality we get  $g(x) \leq c \cdot f(x)$  where  $c = \frac{1}{c_2}$ .
- So then we have  $f(x) \in O(g(x))$  and  $g(x) \in O(f(x))$

2.  $\Leftarrow$

- Suppose  $f(x) \in O(g(x))$  and  $g(x) \in O(f(x))$ .
- Then we know that  $\exists c_1, c_2$  such that  $f(x) \leq c_1 \cdot g(x)$  and  $g(x) \leq c_2 \cdot f(x)$ .
- so  $\frac{1}{c_1}g(x) \leq f(x) \leq c_2 \cdot g(x)$ .
- But this is the definition of  $\Theta$ , therefore  $f(x) \in \Theta(g(x))$ .

- What is the tightest bound we can form here:

1.  $x^2 + 3x + 5 \in \Delta(x^3)$ : big- $O$
2.  $2^n \log(6) + n^2 \in \Delta(2^n)$ :  $\Theta$
3.  $2^n \cdot n! + 2^n \log(n) \in \Delta(2^n)$ :  $\Omega$

- To prove the first one of the previous questions, we use the limit method We have:

$$\lim_{x \rightarrow \infty} \frac{x^2 + 3x + 4}{x^3} = \lim_{x \rightarrow \infty} \frac{2x + 3}{3x^2} = \lim_{x \rightarrow \infty} \frac{2}{6x} = 0$$

Therefore we can conclude that  $x^3$  grows much faster, and so we get that  $x^2 + 3x + 4 \in O(x^3)$

- Next up Sequences: Can we name the first 4 terms of the following sequence  $\{2^n + 1\}_{n=0}^\infty$ ?

1.  $a_0 = 2^0 + 1 = 2$

2.  $a_1 = 2^1 + 1 = 3$

3.  $a_2 = 2^2 + 1 = 5$

4.  $a_3 = 2^3 + 1 = 9$

- Now compute the following sum  $\sum_{i=1}^5 6$  We have:

$$\sum_{i=1}^5 6 = 6 \sum_{i=1}^5 1 = 6 \cdot (5 - 1 + 1) = 6 \cdot 5 = 30$$

- How about the following geometric  $\sum_{i=1}^8 3 \cdot 2^i$

- We can recognize that this expression is the sum of a geometric progression (i.e., geometric series), which is in general given as  $\sum_{i=0}^n ar^i$ , except that here we start at 1 instead of zero, so we can simply compute it by subtracting the first term from the sum.

- The formula for computing this is  $\frac{a \cdot r^{n+1} - a}{r-1}$ , here  $r = 2$  and  $a = 3$

- Using the formula we can get  $\frac{3 \cdot 2^9 - 3}{1} = 3 \cdot 512 - 3 = 1533$

- However, remember we have to subtract the first term, so  $1533 - 3 \cdot 2^0 = 1530$ .