

## Homework 7

**Assigned on:** Mon March 23, 2009.

**Due:** Friday, April 3, 2009.

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You can choose to do either Section 1 (worth 110 points) or Sections 2 to 9 (with 110 points plus 20 bonus points). It is your decision. If you choose to do both options, you will receive the *maximum* of the grades of the two sections, not the sum.

### 1 Implementation: Solving SAT 110 points

Write a search algorithm to determine the satisfiability of a SAT instance. You can either write:

- A DPLL procedure (backtrack search),
- A local search procedure.

You must

- Clearly describe, in addition to your code, your data structures, how your search algorithm operates, and the improvements, if any, that you have included in your code.
- We recommend that you use the standard file for input files known as the ‘simplified version of the DIMACS format’:  
<http://www.satcompetition.org/2009/format-benchmarks2009.html>
- Test the performance of your algorithm on some non trivial uniform random instances taken from the SAT Competition. For example:  
<http://www.cs.ubc.ca/~hoos/SATLIB/benchm.html>

Alert: many implementations exist in the literature and on the web. We expect you to do your *own* implementation.

**2 AIMA, Exercise 6.3, page 190. 10 points**

**3 AIMA, Exercise 7.2, page 236. 16 points**

**4 AIMA, Exercise 7.5, page 237. 6 points**

Note that  $A \Leftrightarrow B \Leftrightarrow C$  is equivalent to  $A \Leftrightarrow B$  and  $B \Leftrightarrow C$ .

## 5 Truth Tables

8 points

Use truth tables to show that each of the following is a tautology.

1.  $(p \wedge q) \rightarrow \neg(\neg p \vee \neg q)$
2.  $[Mary \wedge (Mary \rightarrow Susy)] \rightarrow Susy$
3.  $\alpha \rightarrow [\beta \rightarrow (\alpha \wedge \beta)]$
4.  $(a \rightarrow b) \rightarrow [(b \rightarrow c) \rightarrow (a \rightarrow c)]$

## 6 AIMA, Exercise 7.8, page 237.

16 points

only c, d, e, f, g and h.

## 7 Logical Equivalences

8 points

Using a method of your choice, verify:

1.  $(\alpha \rightarrow \beta) \equiv (\neg\beta \rightarrow \neg\alpha)$  contraposition
2.  $\neg(\alpha \wedge \beta) \equiv (\neg\alpha \vee \neg\beta)$  de Morgan
3.  $(\alpha \wedge (\beta \vee \gamma)) \equiv ((\alpha \wedge \beta) \vee (\alpha \wedge \gamma))$  distributivity of  $\wedge$  over  $\vee$

## 8 Proofs

28 points

Give the explanations of each step if the steps are given, and give both the explanation and step if they are not.

- If  $q \wedge (r \wedge p), t \rightarrow v, v \rightarrow \neg p$ , then  $\neg t \wedge r$ .

**Proof**

**Explanations**

- |                            |       |
|----------------------------|-------|
| 1. $q \wedge (r \wedge p)$ | Given |
| 2. $t \rightarrow v$       | Given |
| 3. $v \rightarrow \neg p$  | Given |
| 4. $t \rightarrow \neg p$  |       |
| 5. $(r \wedge p)$          |       |
| 6. $r$                     |       |
| 7. $p$                     |       |
| 8. $\neg\neg p$            |       |
| 9. $\neg t$                |       |

10.  $\neg t \wedge r$

- If  $p \rightarrow (q \wedge r)$ ,  $q \rightarrow s$ , and  $r \rightarrow t$ , then  $p \rightarrow (s \wedge t)$ .

**Proof**

**Explanations**

- 1.
- 2.
- 3.
- 4.
- 5.
- 6.
- 7.

• **Prove by contradiction.**

If  $\neg(\neg p \wedge q)$ ,  $p \rightarrow (\neg t \vee r)$ ,  $q$ , and  $t$ , then  $r$ .

**Proof**

**Explanations**

1. $\neg(\neg p \wedge q)$	Given
2. $p \rightarrow (\neg t \vee r)$	Given
3. $q$	Given
4. $t$	Given
5. $\neg r$	Negation of Conclusion
6.	
7.	
8.	
9.	
10.	
11.	
12.	

**9 AIMA, Exercise 7.11, page 238. 18 points + 20 bonus**

Parts a, b, and c are required. Parts d, e, and f are bonus.