

Recitation 12.

Section 3.2 & 2.4

P91

3.2: 9: Show that $x^2 + 4x + 17$ is $O(x^3)$ but x^3 is not $O(x^2 + 4x + 17)$.

proof:

① We first proof $x^2 + 4x + 17$ is $O(x^3)$

We know that $x^2 + 4x + 17 < x^3 + x^3 + x^3$ for $x \geq 3$

$$\text{so: } x^2 + 4x + 17 < 3x^3$$

thus for all sufficiently large x

$$x^2 + 4x + 17 < Cx^3 \quad [C=3, x_0=3]$$

which implies that $x^2 + 4x + 17$ is $O(x^3)$

② Next, we proof that x^3 is not $O(x^2 + 4x + 17)$.

We proof it by contradiction.

We assume that x^3 is $O(x^2 + 4x + 17)$.

$$\text{then } x^3 \leq C(x^2 + 4x + 17) \leq C(3x^2)$$

$$\Rightarrow x^3 \leq 3Cx^2$$

$$\text{for } x \neq 0 \Rightarrow x \leq 3C \quad \text{for some } C \text{ and all sufficiently large } x.$$

which is impossible! then we see a contradiction.

Therefore, x^3 is not $O(x^2 + 4x + 17)$.

① For each of the following functions, give and prove a tight asymptotic inclusion of the form $f(n) \in \Delta(g(n))$

Where Δ is one of O, Ω, Θ . ⁽⁴⁾ $f(n) = \log n$ $g(n) = (n + \log n)(n - 1)$

$$\begin{aligned} \text{proof: } & \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \\ &= \lim_{n \rightarrow \infty} \frac{\log n}{(n + \log n)(n - 1)} \\ &= \lim_{n \rightarrow \infty} \frac{\log n}{n^2 + n \log n - n - \log n} \\ &= \lim_{n \rightarrow \infty} \frac{\log n}{n^2 + (n - 1) \log n - n} \\ &= \lim_{n \rightarrow \infty} \frac{1}{\frac{n^2}{\log n} + (n - 1) - \frac{n}{\log n}} \\ &= 0 \end{aligned}$$

$\therefore f(n) \in O(g(n))$

$$(b) f(n) = 2^n, g(n) = 2^{\log n}$$

$$\begin{aligned} \text{proof: } \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} &= \lim_{n \rightarrow \infty} \frac{2^n}{2^{\log n}} \\ &= \lim_{n \rightarrow \infty} 2^{(n - \log n)} \\ &= \infty \end{aligned}$$

$$\therefore f(n) \in \Omega(g(n))$$

(c) Prove that for any real numbers, α, β , such that

$$\alpha > \beta > 0, \quad \beta^n \in O(\alpha^n)$$

$$\begin{aligned} \text{proof: } \lim_{n \rightarrow \infty} \frac{\beta^n}{\alpha^n} &= \lim_{n \rightarrow \infty} \left(\frac{\beta}{\alpha}\right)^n \end{aligned}$$

$$\because \alpha > \beta \quad \therefore 0 < \frac{\beta}{\alpha} < 1$$

$$\therefore \lim_{n \rightarrow \infty} \left(\frac{\beta}{\alpha}\right)^n = 0$$

$$\therefore \beta^n \in O(\alpha^n)$$

●② For the following functions, prove a tight inclusion of the form $f \in \Delta(g)$, where Δ is one of $\mathcal{O}, \Omega, \Theta$

$$(9) \quad f(n) = n^2 \log n \quad g(n) = \left(\sum_{i=1}^n i \right) - \frac{n-1}{2}$$

Solution: $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)}$

$$= \lim_{n \rightarrow \infty} \frac{n^2 \log n}{\frac{n^2+1}{2}}$$

via L'Hôpital's law and the product rule, we have.

$$\lim_{n \rightarrow \infty} \frac{f'(n)}{g'(n)}$$

$$= \lim_{n \rightarrow \infty} \frac{2n \log n + \frac{n^2}{n \ln 2}}{n}$$

$$= \lim_{n \rightarrow \infty} \frac{2n \log n + \frac{n}{\ln 2}}{n}$$

$$= \lim_{n \rightarrow \infty} 2 \log n + \frac{1}{\ln 2}$$

$$= \infty$$

∴ $f(n) \in \Omega(g(n))$.

$$(b) f(n) = \log^2 n, g(n) = 4n$$

$$\text{Solution: } \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{\log^2 n}{4n}$$

$$= \lim_{n \rightarrow \infty} \frac{2 \log n \cdot \frac{1}{n \ln 2}}{4} \quad [\text{L'Hôpital's product rule}]$$

$$= \lim_{n \rightarrow \infty} \frac{\log n}{2n \ln 2}$$

$$= \lim_{n \rightarrow \infty} \frac{1}{\frac{2n \ln 2}{\log n}} \quad [\text{L'Hôpital's law}]$$

$$= \lim_{n \rightarrow \infty} \frac{1}{2 \ln^2 2 n}$$

$$= 0$$

$$\therefore f(n) \in \mathcal{O}(g(n))$$

$$(c) f(n) = n^2, g(n) = n \log(8^n)$$

$$\text{Solution: } \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{n^2}{n \log 8^n}$$

$$= \lim_{n \rightarrow \infty} \frac{n^2}{n^2 \log 8} \quad [\text{divide by } n^2]$$

$$= \lim_{n \rightarrow \infty} \frac{1}{\log 8}$$

$$= \frac{1}{3}$$

$$\text{So, } f(n) \in \Theta(g(n))$$

③ For the next 4 questions, let

$$\begin{aligned}f(n) &= n^2 \\g(n) &= \log^2(n) \\h(n) &= 2^n\end{aligned}$$

(a) Give a tight asymptotic characterization of $g \circ f$

Solution:
$$\begin{aligned}g \circ f &= g(f(n)) = \log^2(n^2) \\&= [2 \log n]^2 \\&= 4 \log^2 n \\&\in \Theta(\log^2 n)\end{aligned}$$

(b) Give a tight asymptotic characterization of $g \circ h$

Solution:
$$\begin{aligned}g \circ h &= g(h(n)) = \log^2(2^n) \\&= [\log 2^n]^2 \\&= [n]^2\end{aligned}$$

(c) Give a tight asymptotic characterization of $f \circ g$

Solution:
$$\begin{aligned}f \circ g &= f(g(n)) = (\log^2 n)^2 \\&= \log^4 n \\&\in \Theta(\log^4 n)\end{aligned}$$

(d) Give a tight asymptotic characterization of $h \circ g$

Solution:
$$\begin{aligned}h \circ g &= h(g(n)) = 2^{\log^2 n} = n^{\log n} \\&\in \Theta(n^{\log n})\end{aligned}$$

④ For the next 3 questions, let

$$f(n) \in \Theta(n^2)$$

$$g(n) \in \Omega(n^3)$$

$$h(n) \in \mathcal{O}(n \log n)$$

Answer the following questions:

(a) Can we say $f \circ g \in \Theta(n^6)$?

Solution:

No, As a counter example: Let $g(n) = n^4$ and

so $g(n) \in \Omega(n^3)$. Let $f(n) = n^2$.

Then $f \circ g = f(g(n)) = n^8 \notin \Theta(n^6)$

(b) Can we say that $f \circ f \in \Omega(n)$?

Solution:

Yes, The lower bound on each function is Linear, and so the composition will not change a linear lower bound.

(c) Can we say that $g(n) \cdot h(n) \in \mathcal{O}(n^4 \log n)$?

Solution:

No, Let $g(n) = n^4$ and $h(n) = n \log n$ then

$$g(n) \cdot h(n) = n^5 \log n \notin \mathcal{O}(n^4 \log n)$$

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3.2 : 35. Suppose that $f(x)$ is $O(g(x))$ where f and g are increasing and unbounded functions. Show that $\log |f(x)|$ is $O(\log |g(x)|)$.

proof: Because $f(x)$ and $g(x)$ are increasing and unbounded, we can assume $f(x) \geq 1$ and $g(x) \geq 1$ for sufficiently large x .

$\therefore f(x)$ is $O(g(x))$

Thus, there are constants C and k , such that

$$f(x) \leq C g(x) \text{ for } x > k.$$

this implies that $\log f(x) \leq \log [C g(x)]$
 $= \log C + \log g(x)$
 $< 2 \log g(x)$ for large x .
[$C=2$]

Hence, $\log f(x)$ is $O(\log g(x))$

Exe. 36 practice after class.

3.2: Prob. 51. (b) show that $x \log x$ is $o(x^2)$

$$\text{proof: } \lim_{x \rightarrow \infty} \frac{x \log x}{x^2}$$

$$= \lim_{x \rightarrow \infty} \frac{\log x}{x}$$

using L'Hôpital's rule

$$= \lim_{x \rightarrow \infty} \frac{(\log x)'}{x'}$$

$$= \lim_{x \rightarrow \infty} \frac{1}{x \ln 2}$$

$$= 0$$

Then, $x \log x$ is $o(x^2)$

3.2: 51 (d) show that $x^2 + x + 1$ is not $o(x^2)$.

$$\text{proof: } \lim_{x \rightarrow \infty} \frac{x^2 + x + 1}{x^2}$$

$$= \lim_{x \rightarrow \infty} \left(1 + \frac{1}{x} + \frac{1}{x^2} \right)$$

$$= 1$$

$$\neq 0$$

Thus, $x^2 + x + 1$ is not $o(x^2)$.

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2.4:19: Show that $\sum_{j=1}^n (a_j - a_{j-1}) = a_n - a_0$ where $a_0 \dots a_n$ is a sequence of real numbers.

Proof: Let's look at $\sum_{j=1}^n (a_j - a_{j-1})$

$$= \cancel{a_1} - a_0 + \cancel{a_2} - \cancel{a_1} + \cancel{a_3} - \cancel{a_2} + \dots + a_n - \cancel{a_{n-1}}$$

$$= a_n - a_0$$

$$\text{so. } \sum_{j=1}^n (a_j - a_{j-1}) = a_n - a_0$$

Exe: Determine which of the following series converge and which diverge. Justify your answer.

a) $\sum_{n=1}^{\infty} \frac{1}{n(n-1)}$

Solution: ① The series converge.

$$\begin{aligned}\sum_{n=1}^{\infty} \frac{1}{n(n-1)} &= (1 - \frac{1}{2}) + (\frac{1}{2} - \frac{1}{3}) + \dots + (\frac{1}{n-1} - \frac{1}{n}) + \dots \\ &= \lim_{n \rightarrow \infty} (1 - \frac{1}{n}) \\ &= 1\end{aligned}$$

b) $\sum_{n=1}^{\infty} 3^n e^{-n}$

Solution: $\sum_{n=1}^{\infty} 3^n e^{-n}$

$$= \sum_{n=1}^{\infty} \left(\frac{3}{e}\right)^n$$

because $\frac{3}{e} > 1$, $\therefore \left(\frac{3}{e}\right)^n$ diverges to ∞

c) $\sum_{n=1}^{\infty} \frac{n}{10n+17}$

Solution: $a_n = \frac{n}{10n+17}$ in this series.

Let's look at the limit of this sequence.

$$\lim_{n \rightarrow \infty} \frac{n}{10n+17} = \lim_{n \rightarrow \infty} \frac{1}{10 + \frac{17}{n}} = \frac{1}{10}$$

Since the limit of the sequence does not equal zero, the n^{th} term test for divergence tells us that the series must diverge.