Computer Science & Engineering 155E Computer Science I: Systems Engineering Focus

Lecture - Recursion

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Recursion

- Functions in C routinely call other functions
- Example: the main function calls the quadraticRootO1 function, which calls the discriminant function, which calls the sqrt function
- C allows functions to also *call themselves*
- This is known as recursion

Recursive Functions

- Recursive functions are common in mathematics
- Sequences are recursively defined functions
- Recall the interpolation method for computing the square root:

$$x_{i} = \frac{1}{2} \left(x_{i-1} + \frac{n}{x_{i-1}} \right)$$

- Functions defined using the functions in the definition (recurrence relations
- Canonical example: the Fibonacci sequence

Fibonacci Sequence

- Fibonacci sequence defined as the sum of its two previous elements
- Named for Leonardo of Pisa, known as Fibonacci (a contraction of filius Bonaccio, "son of Bonaccio")

$$F_n = \begin{cases} 0 & \text{if } n = 0\\ 1 & \text{if } n = 1\\ F_{n-1} + F_{n-2} & \text{if } n > 1 \end{cases}$$

$$0, 1, 1, 2, 3, 5, 8, 13, 21, 34, \ldots$$

Sequence:

Fibonacci Sequence

C function

- ► We can easily translate a recursive mathematical function to a recursive
- We just have to be careful to handle certain issues
- ▶ We design a function that calls itself: a function that calls a function of the same name

Precursion C Code f = finclude <stdib.h> f = include <stdib.h+ f = include <stdib.h+ f = include <std

Recursion Rules

When using recursion, some rules must be followed:

- $1. \ \mbox{The function} \ \mbox{must} \ \mbox{have at least} \ \mbox{one} \ \ \mbox{terminating condition}$
- 2. The function must make progress toward a terminating condition

Recursion

Terminating Condition

- ▶ We need some guarantee that a recursive function will eventually halt
- ▶ A recursive function must have at least one *terminating condition*
- ► A "base case" in which the function does not call itself again
- ► In the Fibonacci program: Three terminating conditions
- Each returns a specific value without calling fibonacci again

Recursion Tracing a Recursive Call Progress Must, in some way, make progress toward the terminating condition ► To understand recursion, it is helpful to trace a recursive function call Incrementing/decrementing the passed values Example: fibonacci(5): on the first call, the function makes two Fibonacci example: each recursive call, fibonacci(n-1), calls:fibonacci(4) and fibonacci(3) fibonacci(n-2) decrements \boldsymbol{n} Each one makes two recursive calls, and each one of those makes its own recursive calls, etc. Progress is made toward the terminating conditions Full computation can be illustrated with a *recursion tree* \blacktriangleright Out of bounds check: function is undefined for n<0If all possibilities are not handled: infinite recursion



Tracing a Recursive Call

- ▶ Note: fibonacci(5) required 15 calls to fibonacci
- Many calls were unnecessary: fibonacci(2) was called three times!
- Extensive recomputation is required in this case
- Number of recursive calls is *exponentially* large
- Better way of computing F_n?







Advantages & Disadvantages I

Overview

- Some languages may not support recursion
- Non-trivial fact: Any recursive function can be made non-recursive
- Arguments for and against recursion exist

Advantages:

- Simplified code
- Closely matches a Divide & Conquer approach to problems solving

Advantages & Disadvantages II

Disadvantages:

- Generally inefficient: requires many system stack swaps
- May needlessly recompute values (Fibonacci sequence)
- May be harder to debug and/or consider all possibilities

Better Alternatives:

- Tail recursion (no local state to take up the program stack)
- Use smarter data structures (stacks)
- Use memoization (use of a table to store function values to avoid repeating the same call)

Exercise I

Exercise

In class exercise: Write a non-recursive function for the fibonacci sequence. Modify the main driver program to count the number of additions that are preformed (this will require a global variable) and compare the performance of the two functions.

Exercise II I

The binomial coefficients, C(n,k) or $\binom{n}{k}$ ("*n* choose *k*"), are defined as the number of ways you can select *k* distinct items from a collection of *n* items. A direct combinatorial definition is

$$\binom{n}{k} = \frac{n!}{k!(n-k)!}$$

An alternative is Pascal's identity, which gives a recurrence to compute this value:

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$$

Where $\binom{n}{0} = 1$ for any n and for all k > n, $\binom{n}{k} = 0$. Finally, $\binom{n}{1} = n$.

Exercise II II

Exercise

Write a recursive function to compute $\binom{n}{k}$ using this formula. Then write a function that uses the factorial definition and try to compute $\binom{30}{12}$ with each one. What answers do you get and why? Write a main function that takes n, k as command line arguments and outputs the result of $\binom{n}{k}$ for both the recursive definition and for the factorial definition.