Computer Science & Engineering 155E Computer Science I: Systems Engineering Focus

Lecture - Recursion

Christopher M. Bourke cbourke@cse.unl.edu

Recursion

- ► Functions in C routinely call other functions
- ► Example: the main function calls the quadraticRootO1 function, which calls the discriminant function, which calls the sqrt function
- ▶ C allows functions to also *call themselves*
- ► This is known as recursion

Recursive Functions

- ▶ Recursive functions are common in mathematics
- Sequences are recursively defined functions
- ▶ Recall the interpolation method for computing the square root:

$$x_{i} = \frac{1}{2} \left(x_{i-1} + \frac{n}{x_{i-1}} \right)$$

- Functions defined using the functions in the definition (recurrence relations
- ► Canonical example: the Fibonacci sequence

Fibonacci Sequence

- ▶ Fibonacci sequence defined as the sum of its two previous elements
- Named for Leonardo of Pisa, known as Fibonacci (a contraction of filius Bonaccio, "son of Bonaccio")

$$F_n = \begin{cases} 0 & \text{if } n = 0\\ 1 & \text{if } n = 1\\ F_{n-1} + F_{n-2} & \text{if } n > 1 \end{cases}$$

Sequence:

 $0, 1, 1, 2, 3, 5, 8, 13, 21, 34, \dots$

Fibonacci Sequence

From Math to C

- We can easily translate a recursive mathematical function to a recursive C function
- ▶ We just have to be careful to handle certain issues
- ► We design a function that calls itself: a function that calls a function of the same name

Recursion

Recursion

Rules

When using recursion, some rules must be followed:

- 1. The function must have at least one terminating condition
- 2. The function must make progress toward a terminating condition

Recursion

Terminating Condition

- ▶ We need some guarantee that a recursive function will eventually halt
- ▶ A recursive function must have at least one *terminating condition*
- ▶ A "base case" in which the function does not call itself again
- ▶ In the Fibonacci program: Three terminating conditions
- ► Each returns a specific value without calling fibonacci again

Recursion

Progress

- ▶ Must, in some way, make progress toward the terminating condition
- ▶ Incrementing/decrementing the passed values
- ▶ Fibonacci example: each recursive call, fibonacci(n-1), fibonacci(n-2) decrements n
- ▶ Progress is made toward the terminating conditions
- lackbox Out of bounds check: function is undefined for n<0
- ▶ If all possibilities are not handled: infinite recursion

Tracing a Recursive Call

- ▶ To understand recursion, it is helpful to trace a recursive function call
- ► Example: fibonacci(5): on the first call, the function makes two calls:fibonacci(4) and fibonacci(3)
- ► Each one makes two recursive calls, and each one of those makes its own recursive calls, etc.
- ▶ Full computation can be illustrated with a recursion tree

Tracing a Recursive Call

Tracing a Recursive Call

- ▶ Note: fibonacci(5) required 15 calls to fibonacci
- ▶ Many calls were unnecessary: fibonacci(2) was called three times!
- ▶ Extensive recomputation is required in this case
- ▶ Number of recursive calls is *exponentially* large
- ▶ Better way of computing F_n ?

Factorial Example

► Recall the factorial function:

$$n! = n \times (n-1) \times (n-2) \times \cdots, 2 \times 1$$

▶ We can also write a recursive function for this function:

$$F_n = \left\{ \begin{array}{ll} 1 & \text{if } n=1 \\ F_{n-1} \times n & \text{otherwise} \end{array} \right.$$

- Strategy:
 - 1. Identify and handle the base case(s)
 - 2. Identify and handle the recursive call

Factorial C Code

```
finclude<stdlib.h>
finclude<stdlib.h>
finclude<stdlib.h>

int factorial(int n);

fint main(int argc, char *argv[])

fint fi(argc != 2)

fint fi(argc != 2)

fint exit(-1);

fint n = atoi(argv[1]);

fint n = atoi(argv[1]);

fint factorial(int n)

fint factorial(int n);

fint factorial(
```

Factorial Recursion Tree



Factorial Example

- ▶ 5! only requires five calls to factorial
- ▶ Recursion is linear in depth and number of calls

Advantages & Disadvantages I

Overview

- ▶ Some languages may not support recursion
- ▶ Non-trivial fact: Any recursive function can be made non-recursive
- Arguments for and against recursion exist

Advantages:

- ► Simplified code
- ▶ Closely matches a Divide & Conquer approach to problems solving

Advantages & Disadvantages II

Disadvantages:

- ▶ Generally inefficient: requires many system stack swaps
- ► May needlessly recompute values (Fibonacci sequence)
- ▶ May be harder to debug and/or consider all possibilities

Better Alternatives:

- ▶ Tail recursion (no local state to take up the program stack)
- ▶ Use smarter data structures (stacks)
- Use memoization (use of a table to store function values to avoid repeating the same call)

Exercise I

Exercise

In class exercise: Write a non-recursive function for the fibonacci sequence. Modify the main driver program to count the number of additions that are preformed (this will require a global variable) and compare the performance of the two functions.

Exercise II II

Exercise

Write a recursive function to compute $\binom{n}{k}$ using this formula. Then write a function that uses the factorial definition and try to compute $\binom{30}{12}$ with each one. What answers do you get and why? Write a main function that takes n,k as command line arguments and outputs the result of $\binom{n}{k}$ for both the recursive definition and for the factorial definition.

Exercise II I

The binomial coefficients, C(n,k) or $\binom{n}{k}$ ("n choose k"), are defined as the number of ways you can select k distinct items from a collection of n items. A direct combinatorial definition is

$$\binom{n}{k} = \frac{n!}{k!(n-k)!}$$

An alternative is Pascal's identity, which gives a recurrence to compute this value:

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$$

Where $\binom{n}{0}=1$ for any n and for all k>n, $\binom{n}{k}=0$. Finally, $\binom{n}{1}=n$.